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# LOCAL DECODERS AND THRESHOLDS OF TOPOLOGICAL QUANTUM CODES

Aleksander Kubica

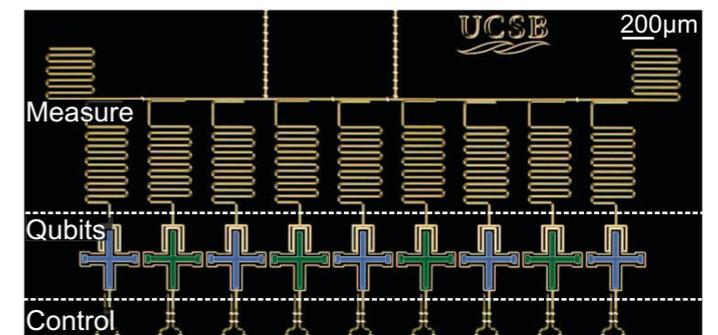
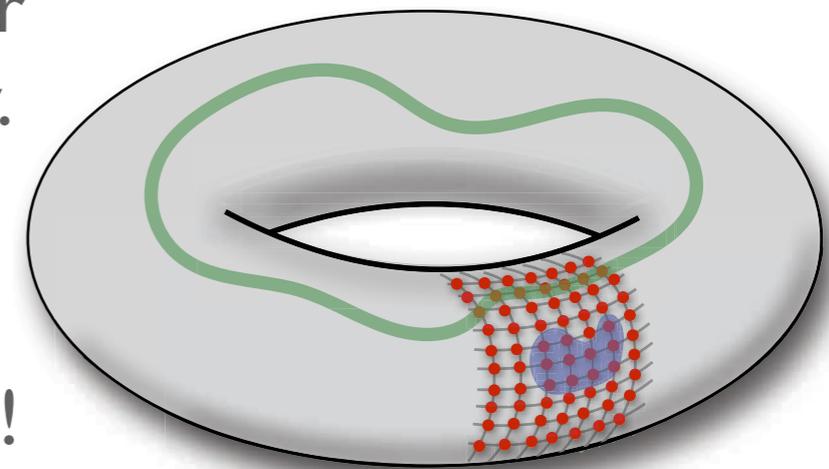


work w/ M. Beverland, F. Brandão, N. Delfosse, J. Preskill, K. Svore

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# TOPOLOGICAL STABILIZER CODES

- Qubits on a manifold, (geometrically) local stabilizer generators, logical information encoded non-locally.
- Well-known models: toric and color codes.
- Can be built in the lab: 2D and local measurements!
- Desired properties:
  - fault-tolerant logical gates,
  - efficient decoders,
  - high threshold.
- **Decoder:** algorithm finding correction from stabilizer measurements.
- **Threshold**  $p_{th}$  = max error rate the code & decoder can tolerate.



Kelly et al., Nature 519, (2015)

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# OUTLINE

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This talk: local (in space/time) decoders w/ provable thresholds.

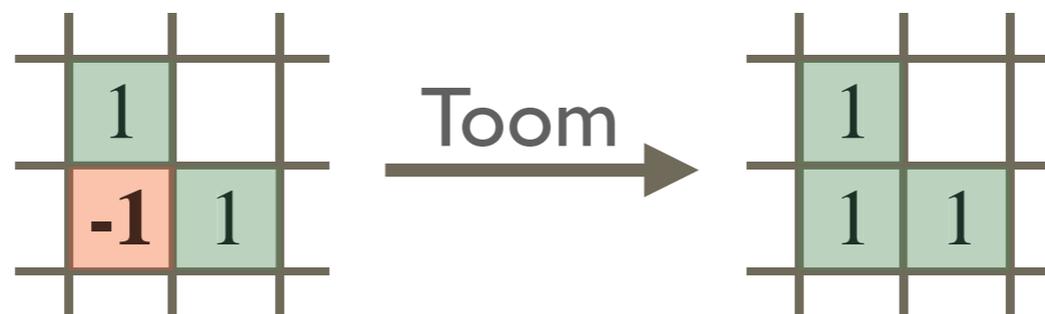
Many toric/color code decoders: non-local, local but heuristic (*Harrington, Dennis, Fowler, Breuckmann, Herold, Duclos-Cianci, Haah, Hastings, Brown,...*).

1. Generalization of Toom's rule to any lattice.
2. Local TC decoder w/ non-zero threshold.
3. Reduction of CC decoding to TC decoding.
4. 3D CC thresholds via stat-mech mappings (arXiv: **1708.07131**)

# NEED FOR (LOCAL) ERROR CORRECTION

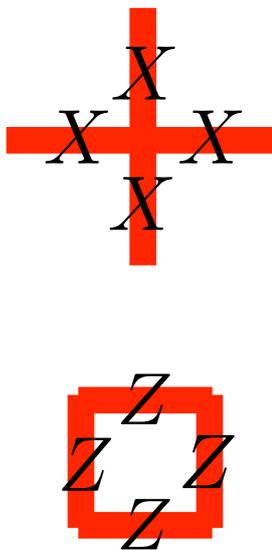
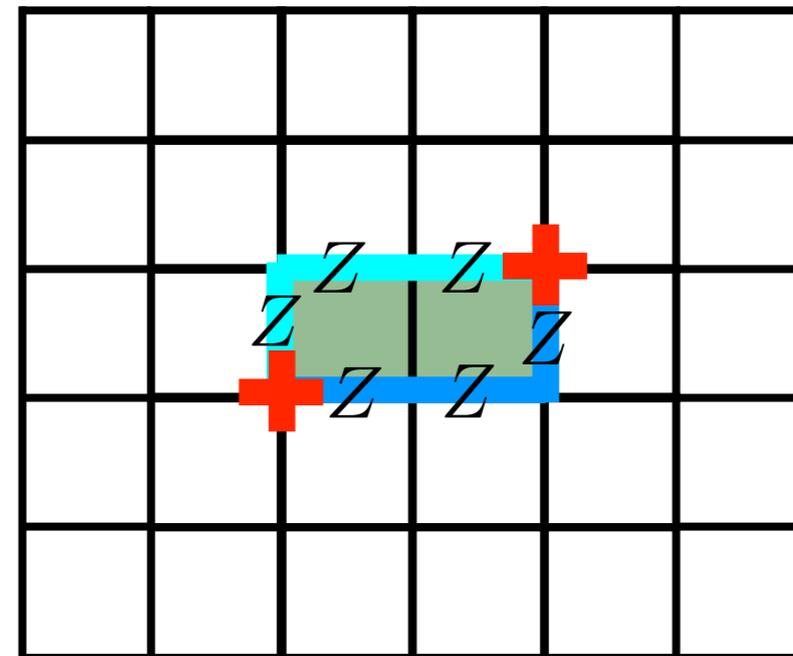
- Errors can accumulate! To prevent that — diagnose and correct errors.
- **Example:** classical memory protecting one bit  $\pm 1$ 
  - repetition code,
  - decoder — majority vote.
- Noise flips some bits. Collecting (global) information takes time — new errors can appear!
- **Goal:** suppress/remove errors by local operations.
- **Toom's rule:** flip bit (face) if it differs from both N and E neighbors.

1	1	1	1	1
1	-1	-1	1	1
1	1	-1	1	1
1	1	1	1	1
-1	1	1	1	-1



# DECODING PROBLEM

- Unlike classical bits, quantum information can't be accessed directly.
- **Stabilizer (CSS) codes:** measure  $X/Z$ -stabilizers and correct  $Z/X$ -errors separately. We consider ideal measurements.
- **Decoding:** find position of errors from violated stabilizers (excitations).
- **2D toric code (Kitaev):**
  - qubits = edges,
  - stabilizers =  $Z$ -faces &  $X$ -vertices,
  - $Z$ -errors = edges,
  - excitations = vertices.
- Decoding successful if error and correction differ by stabilizer.

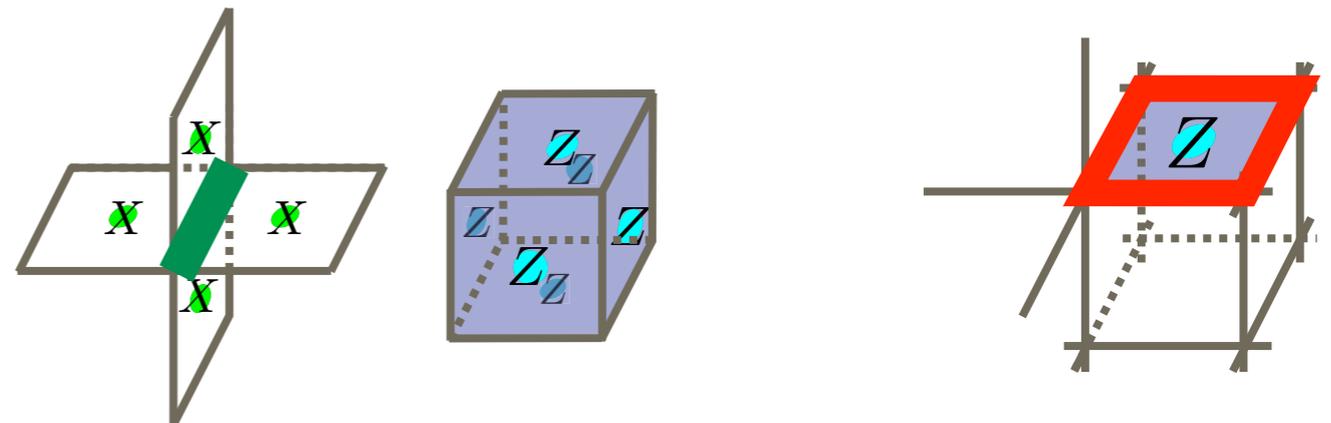


# TOOM'S RULE AS DECODER

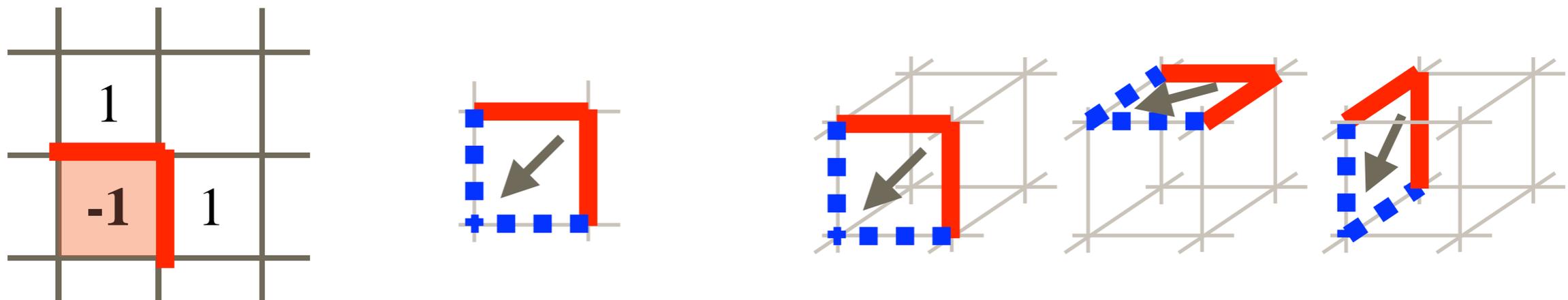
- Toric/color code in  $d$  dim w/  $(k-1)$ - &  $(d-k-1)$ -dim excitations,  $k=1, \dots, d-1$ .

- **3D toric code:**

- qubits = faces,
- stabilizers =  $X$ -edges,  $Z$ -cubes.



- Decode  $Z$  errors = flip faces w/ boundary matching loop-like excitations.
- Toom's rule — a rule for (re)moving domain walls, i.e. “move NE corners”.





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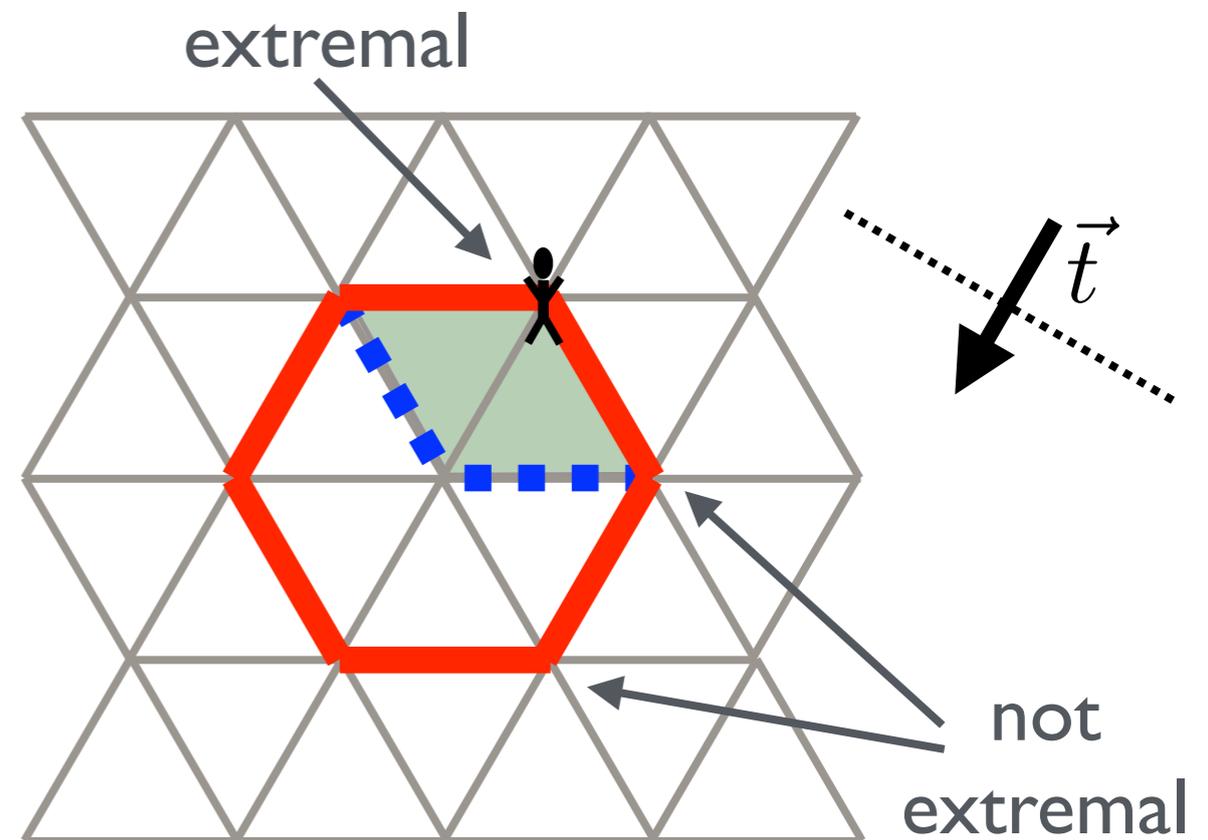
# LOCAL EFFICIENT DECODERS: TORIC AND COLOR CODES

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- **Questions:**
  - Toom's rule on any lattice?
  - does decoding w/ Toom's rule work?
- **Sweep Rule** — a generalization of Toom's rule to any  $d$ -dim lattice and  $k$ -dim domain walls for  $k=1, \dots, d-1$ .
- **Threshold** for local toric code decoders based on the Sweep Rule.
- **Local color code decoders** in  $d \geq 3$  dim by using any toric code decoder.

# SWEEP RULE

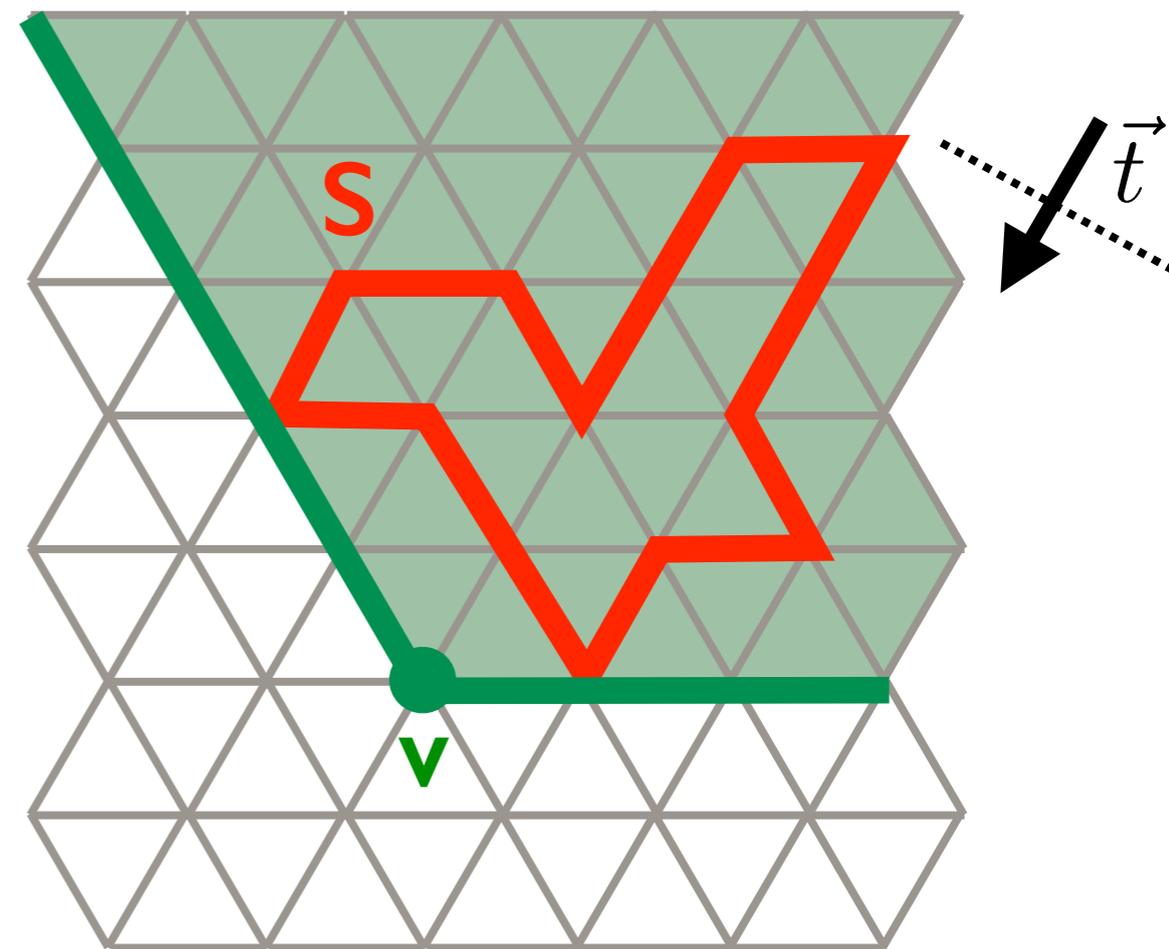
- Change of perspective: not faces but vertices!
- Introduce the sweep direction.
- **Extremal vertex  $v$ :** local restriction of the domain wall is in the sweep direction from  $v$ .



- **Sweep Rule:** if vertex extremal, flip faces in the sweep direction.
- Sweep Rule in  $d \geq 2$  dim defined similarly. Important to flip right cells!

# PROPERTIES OF SWEEP RULE

- The sweep direction induces a partial order  $\leq$  over the set of vertices. Alternative picture: vertices in spacetime, path & causal path.
- Two notions:
  - $\text{cone}(v) = \{\text{vertices } u \mid u \leq v\}$
  - $\text{sup}(S) = \text{least lower bound of } S$
- **Correction region:**  
domain wall  $S$  stays within  $\text{cone}(\text{sup}(S))$ .
- **Monotone:**  
max length of the causal path between  $\text{sup}(S)$  and any vertex of domain wall  $S$ .



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# SWEEP DECODER

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- Toric code in  $d \geq 3$  dim w/  $k$ -dim excitations for  $k=1, \dots, d-2$ .
- **Sweep Decoder**
  1. repeat  $M$  times: simultaneously apply the Sweep Rule for every vertex  $v$ ,
  2. correction = flipped faces.
- Decoder can fail because:
  - (a) domain walls have not been removed in  $M$  time steps,
  - (b) correction introduced logical error.
- Our result: Sweep Decoder has **non-zero threshold**  $p_c$ :  
if error rate  $p \leq p_c$  then  $\text{pr}(\text{success}) \rightarrow 1$  in the limit of lattice size  $L \rightarrow \infty$ .

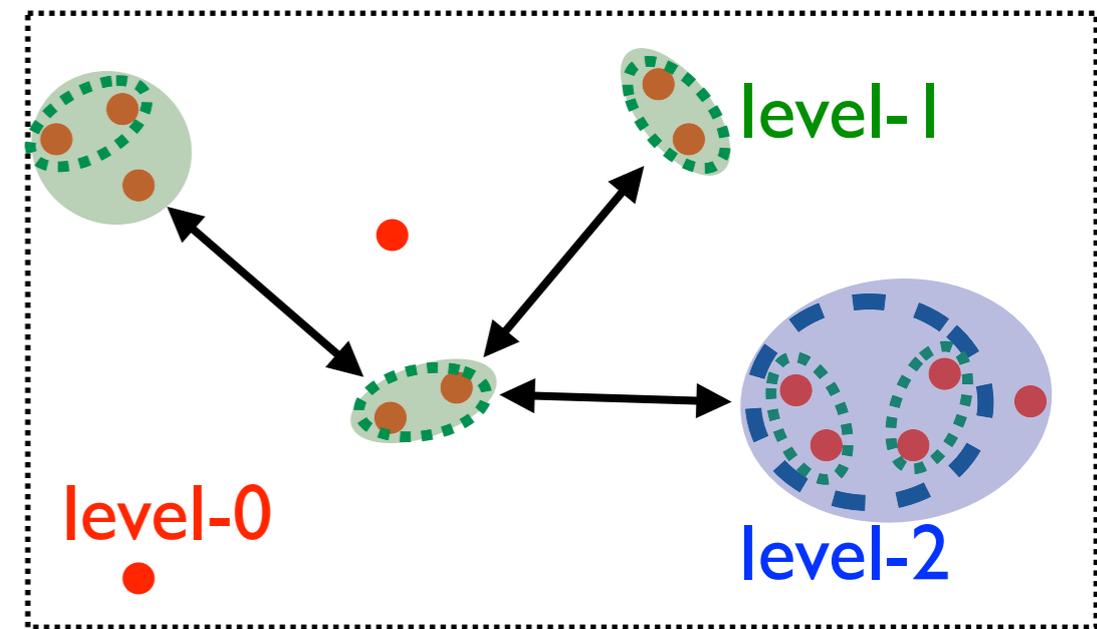
# KEY LEMMAS

- We use ideas of Gacs, Harrington, Bravyi&Haah.
- Level-0 chunk = single error, level-1 chunk = nearby pair of errors, ..., level-n chunk = two disjoint level-(n-1) chunks & diameter  $\leq Q^n/2$ .

- **Lemma 1:** for sufficiently small  $p$  the probability of having a level-n chunk is suppressed doubly exponentially in  $n$ .

- Level-n error  $E_n$  = union of level-n chunks.

- Disjoint decomposition of errors:  
 $E = (E_0 - E_1) + (E_1 - E_2) + \dots + (E_{m-1} - E_m) + E_m$ .



- **Lemma 2:** if  $C$  is a level- $i$  cluster of errors in  $(E_i - E_{i+1})$ , then  $C$  is “not too big” ( $\text{diam}(C) \leq Q^i$ ) and “far from other errors” ( $d(M, (E_i - M)) \geq Q^{i+1}/3$ )<sub>12</sub>

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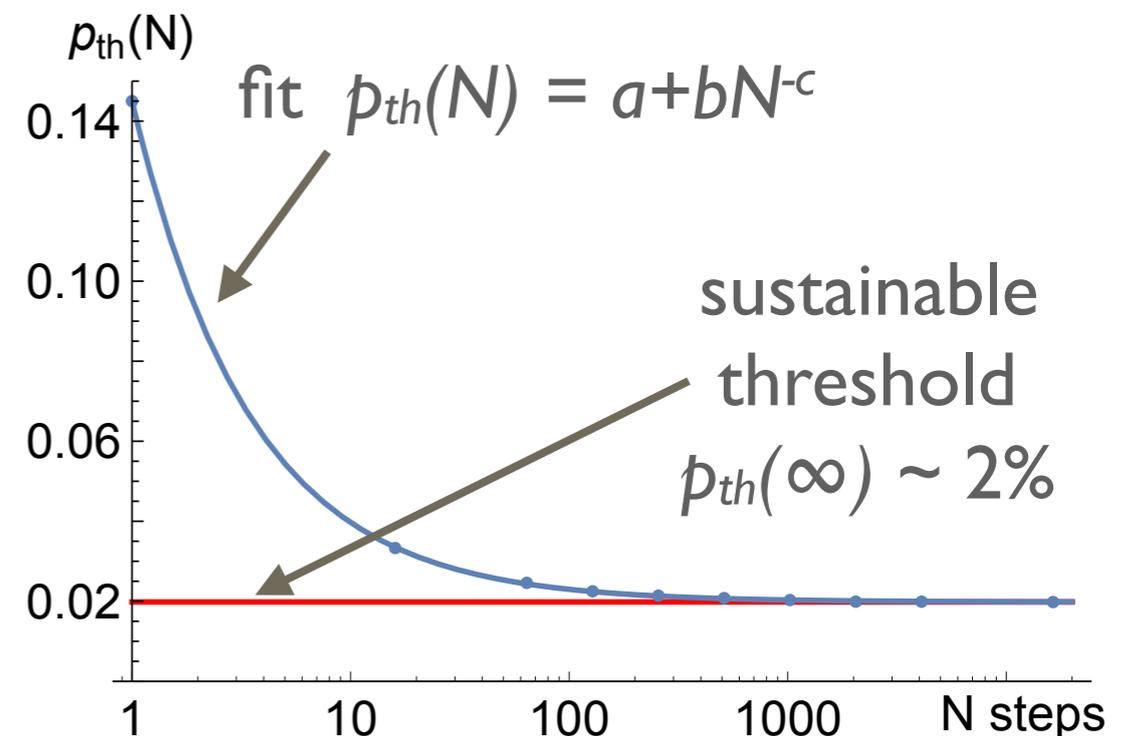
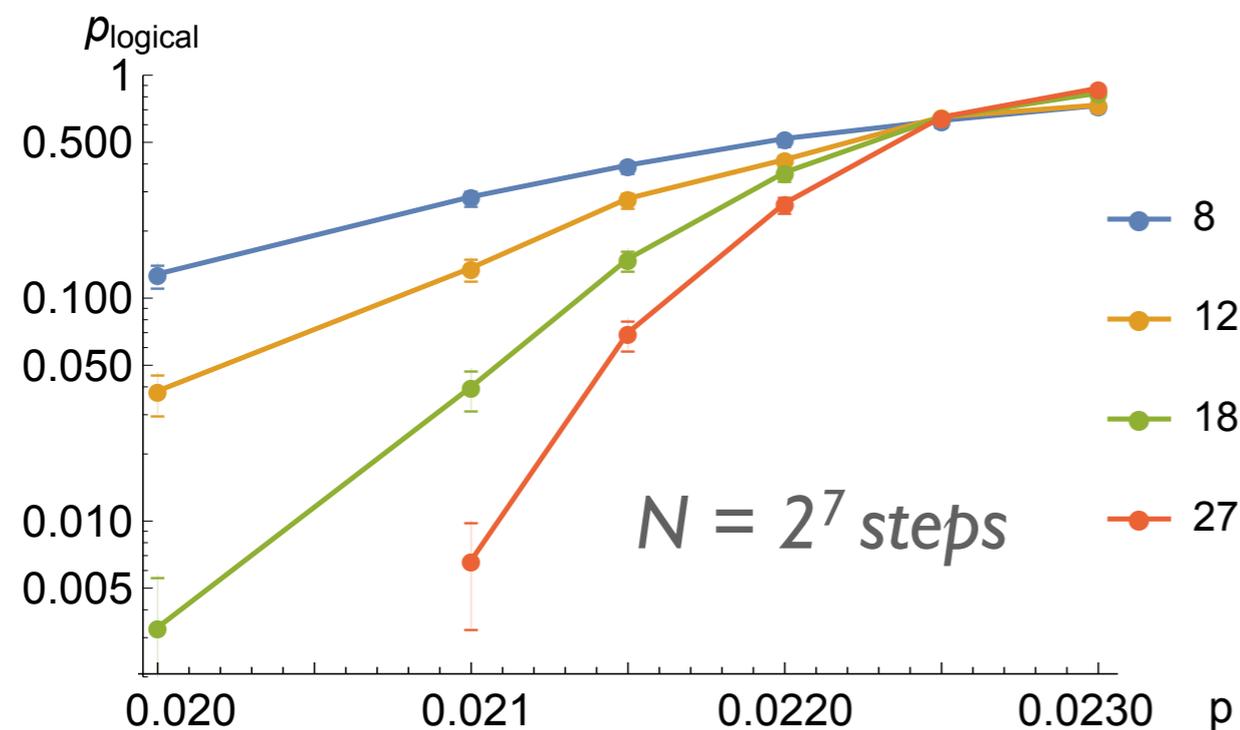
# PUTTING THINGS TOGETHER

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- Assumptions on the lattice: (locally) Euclidean, ...
- Isolated error removed in 1 step, level-1 cluster in  $\sim Q$  steps, ..., level- $i$  cluster  $C$  in time  $\sim Q^i$  (use:  $C$  “not too big” & monotone).
- Removal of level- $i$  cluster  $C$  unaffected by other level- $j$  clusters for all  $j \geq i$  (use:  $C$  “far from other errors” & correction region).
- Correction of  $C$  inside the cone of its boundary, which for low-level clusters ( $i \leq \log L$ ) is a correctible region — no logical error!
- Run local updates for total time  $\sim Q^i$ , where  $i \sim \log L$ . Higher-level clusters might not be removed, but they are very unlikely!
- Failure of the decoder due to presence of high-level clusters:  
 $\text{pr}(\text{fail}) \leq \text{poly}(L) \exp(-cL) \rightarrow 0$  as  $L \rightarrow \infty$ .

# NOISY MEASUREMENTS: 3D TORIC CODE NUMERICS

- Realistic setting — noisy measurements w/ prob =  $p$ .
- Iterate  $N$  times: add new errors, imperfectly measure stabilizers, apply one round of local correction everywhere.
- After  $N$  iterations: measure perfectly, decode, check for logical errors.
- We can find threshold  $p_{th}(N)$  and analyze its behavior in the limit  $N \rightarrow \infty$ .



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# OUTLINE

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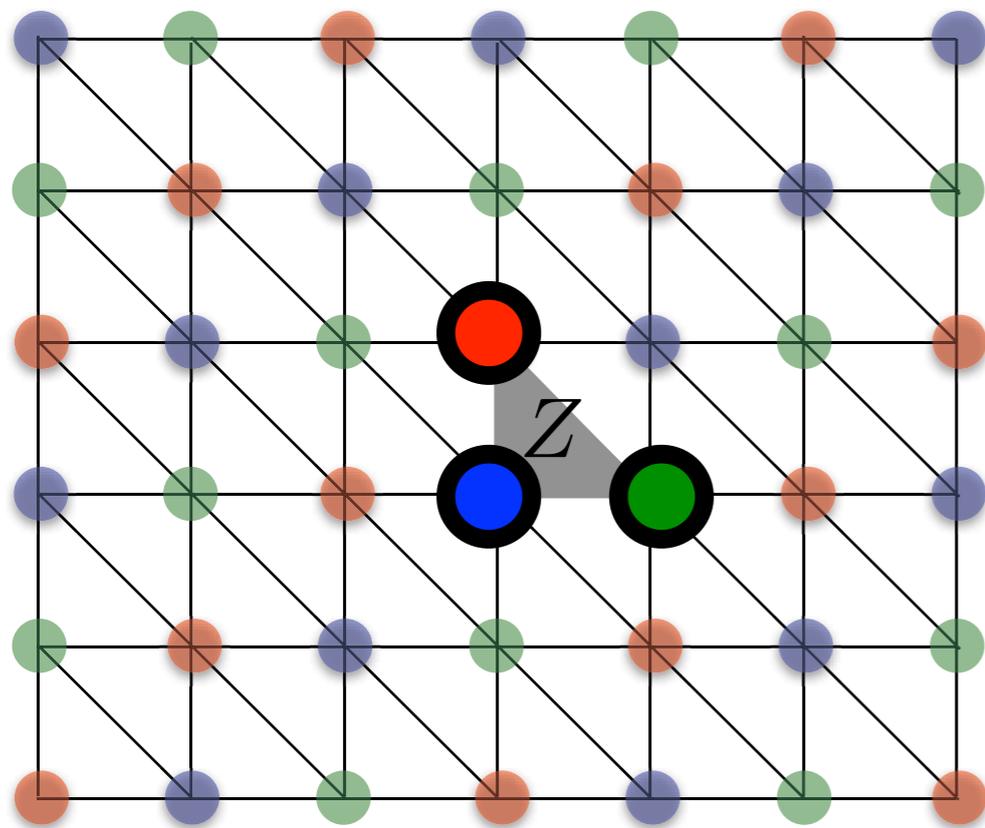
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# TOPOLOGICAL CODE: 2D COLOR CODE

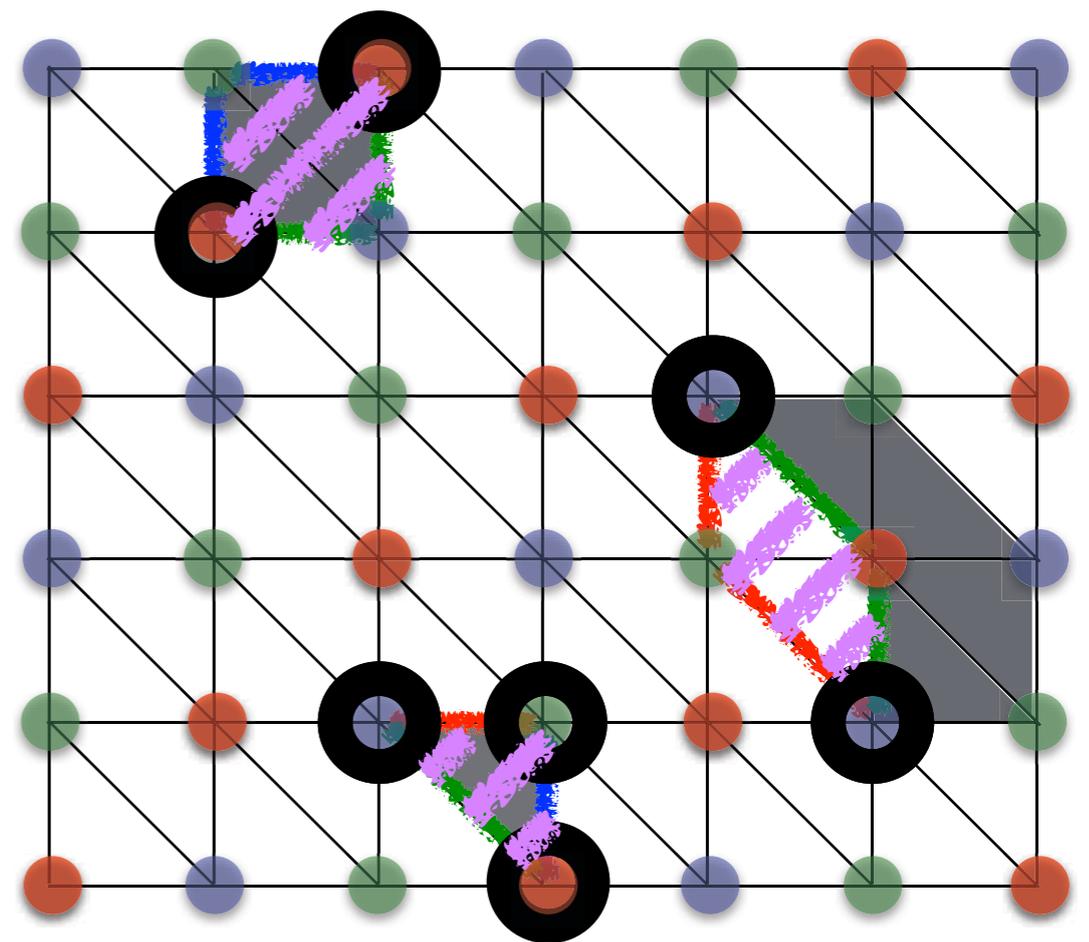
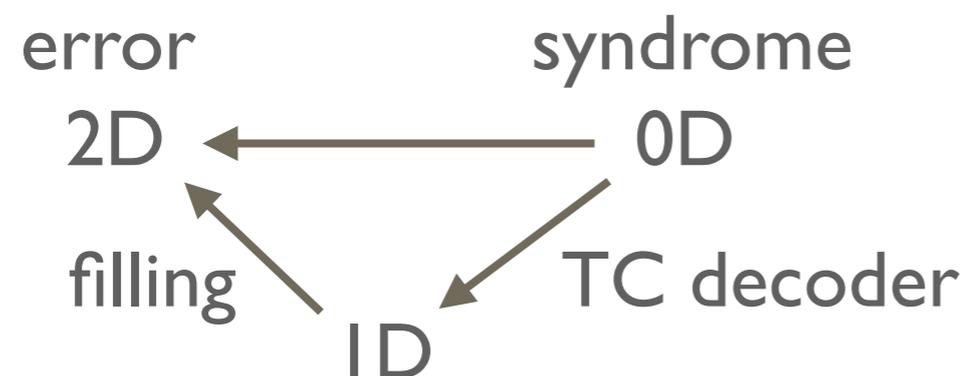


- (Dual) lattice: made of triangles and vertices are 3-colorable.
- 2D color code (*Bombin*):
  - qubits = triangles,
  - stabilizers = X- & Z-vertices.
- Logical Clifford gates are transversal, code switching and gauge fixing, ...
- Decoding seems to be more challenging: excitations created in triples!



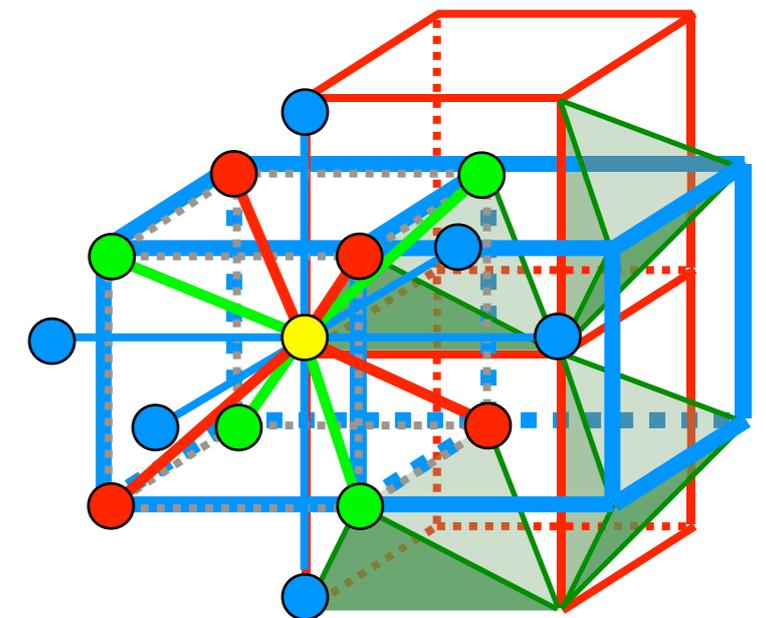
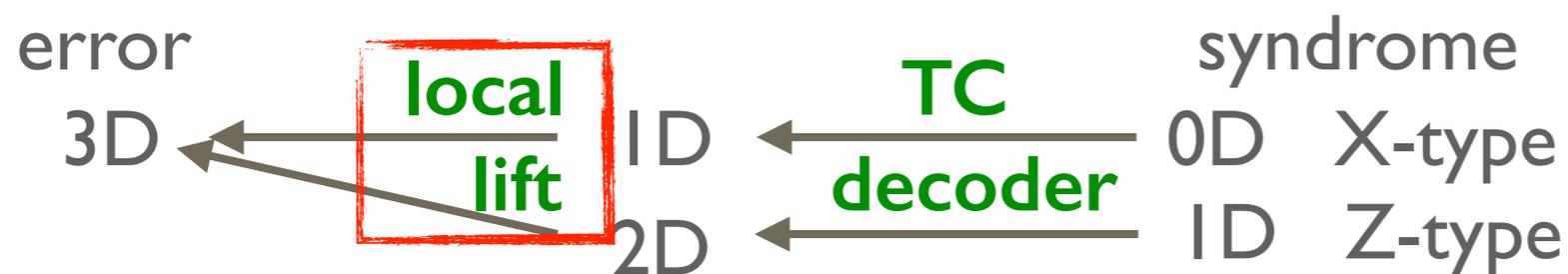
# HOW TO DECODE COLOR CODES?

- **Idea:** color and toric codes are related (*Kubica et al.'15*) — can we use existing toric code decoders?
- Noise changes — correlated errors!
- 2D projection decoder (*Delfosse'14*)
  - TC decoder on three sublattices,
  - global filling.



# LOCAL COLOR CODE DECODERS IN $D \geq 2$ DIM

- Beyond 2D not really explored! Similar ideas work.
- Our result: decoder w/ local reduction and lifting in  $d \geq 2$  dim.
- **3D color code (bcc lattice):**
  - qubits = tetrahedra,
  - stabilizers = X-vertices and Z-edges.



- Any toric code decoder can be used! Fully local for loop-like excitations.
- Toric code thresholds allows to lower bound color code threshold!

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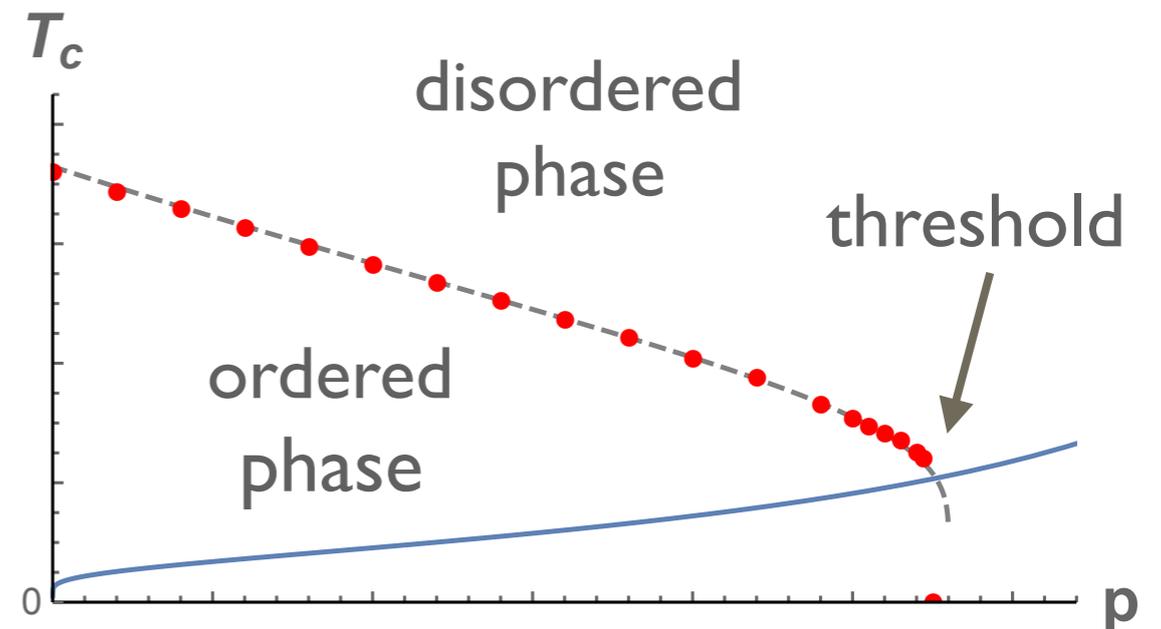
# OUTLINE

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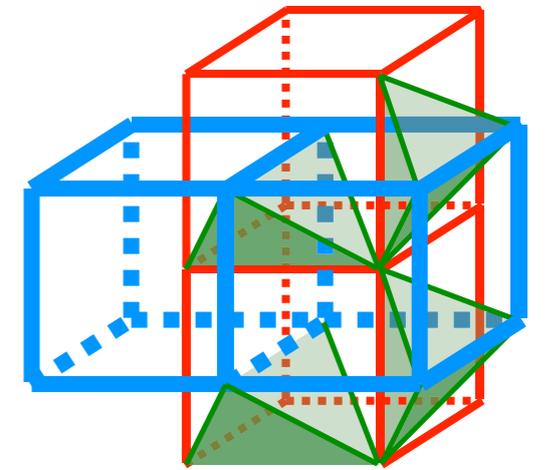
# THRESHOLDS FROM STATISTICAL MECHANICS

- Analytic bounds on threshold (very) low and far from actual values.
- Values of thresholds relevant for: overhead estimates, comparing codes and decoders, experiment, ...
- *Dennis et al.'02*: connection between toric code decoding and a classical spin model (random-bond Ising).
- Ordered phase = successful correction.
- Our results: new spin models relevant for 3D color code & their phase diagrams, thresholds of 3D color code.



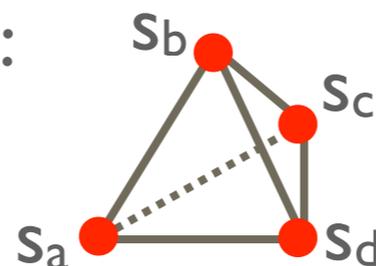
# RANDOM COUPLING ISING MODEL AND 3D COLOR CODE

- 3D bcc lattice:
  - qubits = tetrahedra,
  - stabilizers = X-vertices (A) and Z-edges (B).
- Z/X-errors lead to 0D point-/ 1D loop-like excitations. Logical Z/X operators are 1D string-/ 2D sheet-like.



- 4-body w/ spins on vertices (A):

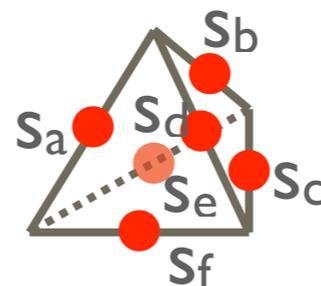
$$H = - \sum \mathcal{K}_{abcd} s_a s_b s_c s_d$$



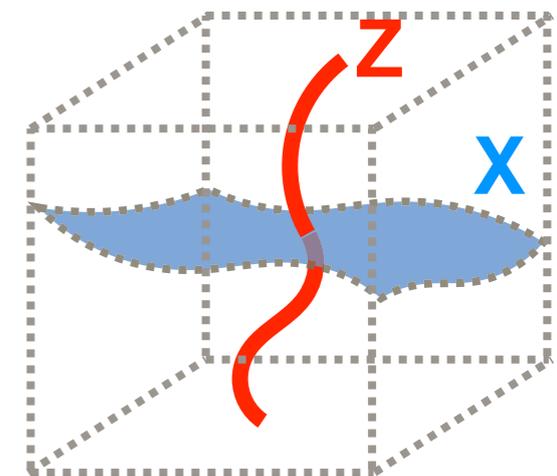
$$\mathcal{K}_{abcd} = \pm 1$$

- 6-body w/ spins on edges (B):

$$H = - \sum \mathcal{K}_{abcdef} s_a s_b s_c s_d s_e s_f$$

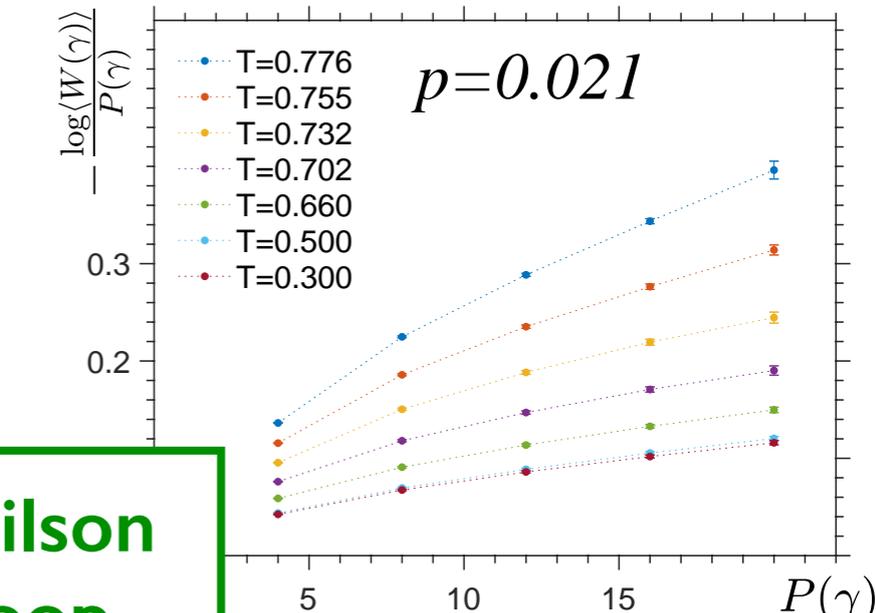
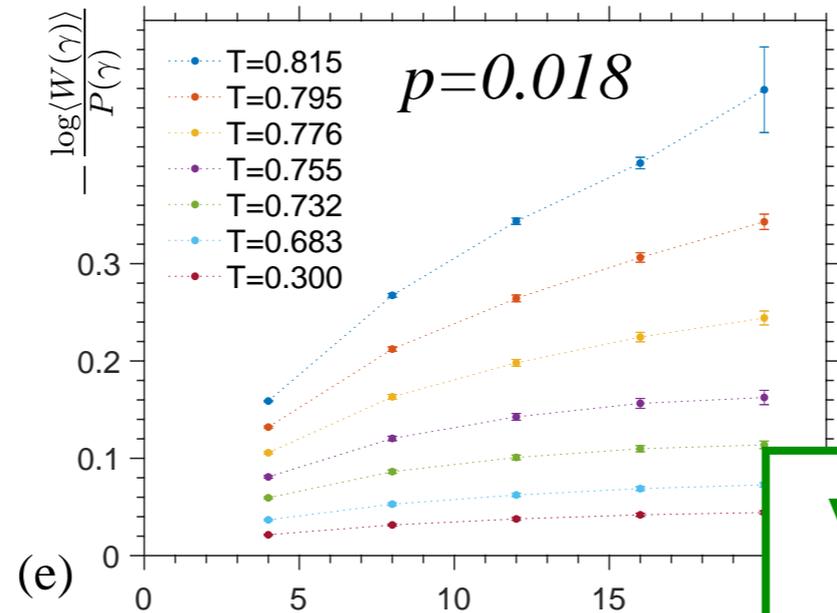
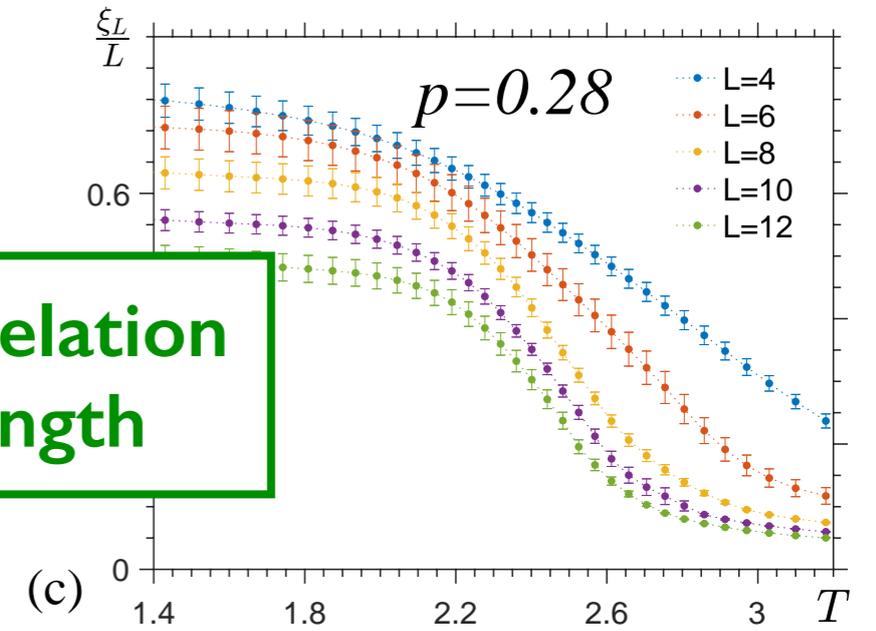
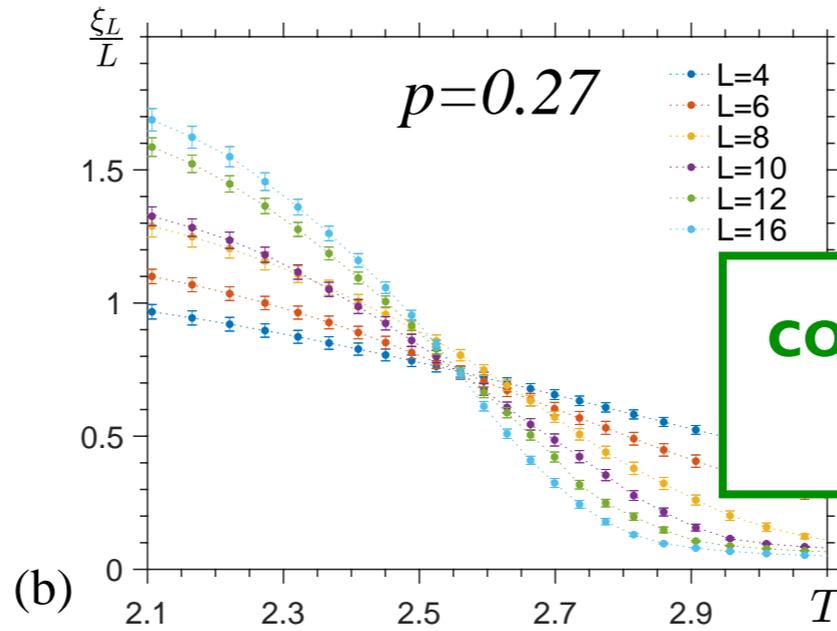
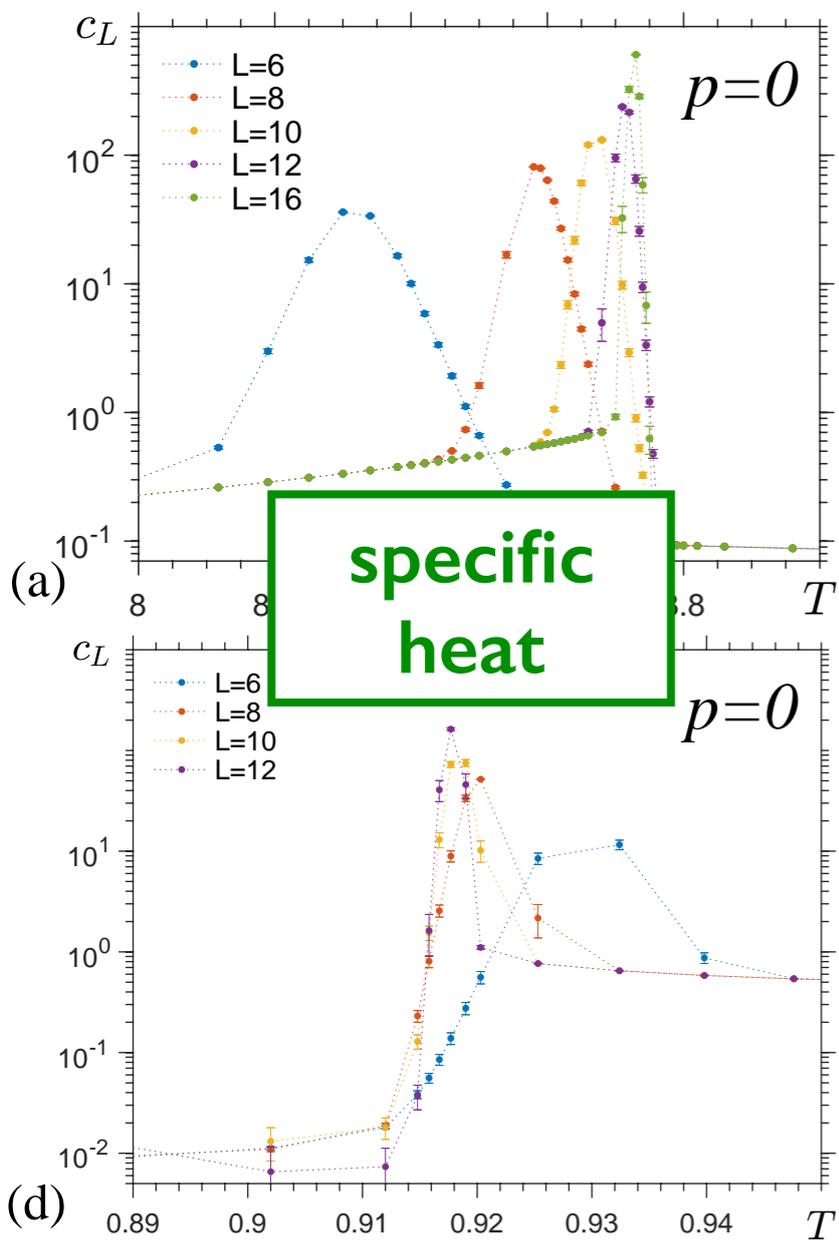


$$\mathcal{K}_{abcdef} = \pm 1$$

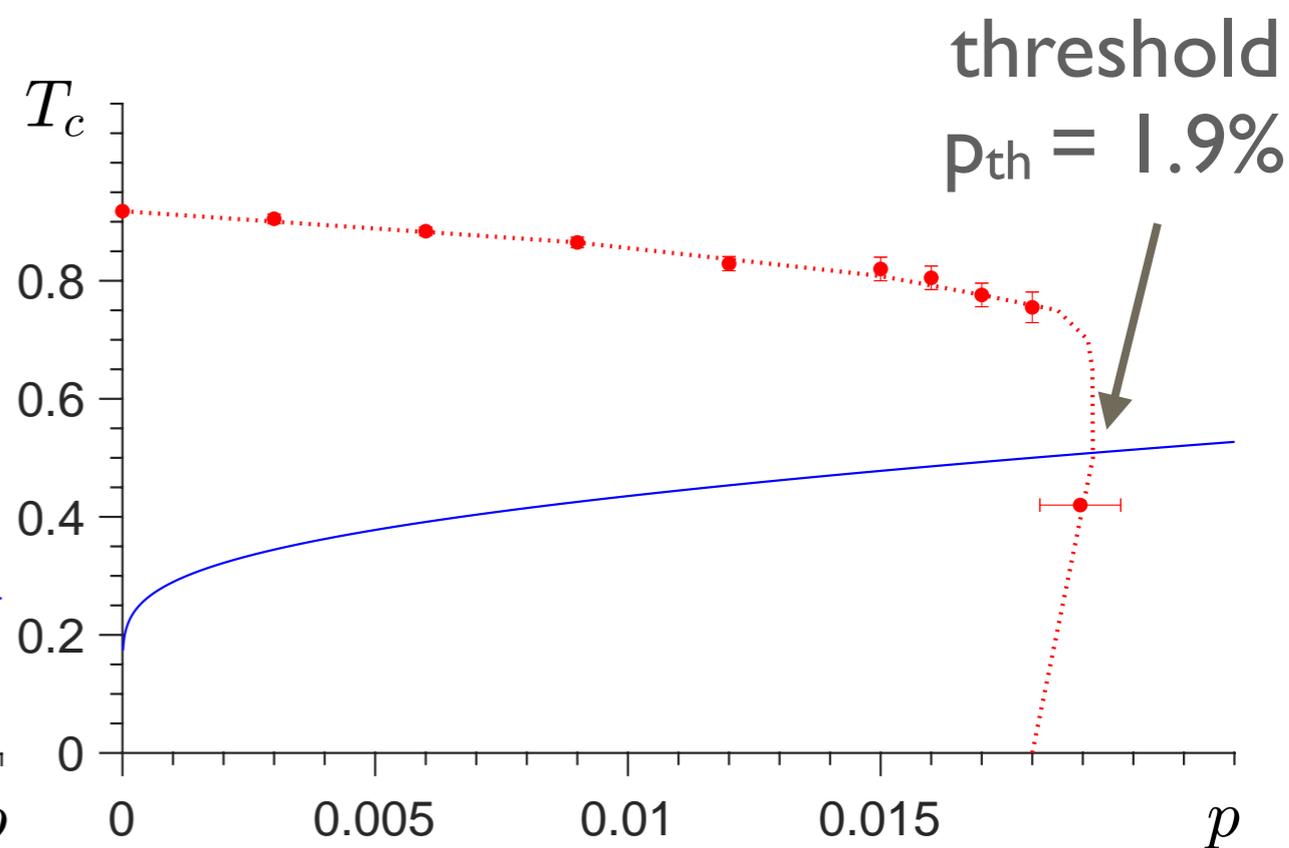
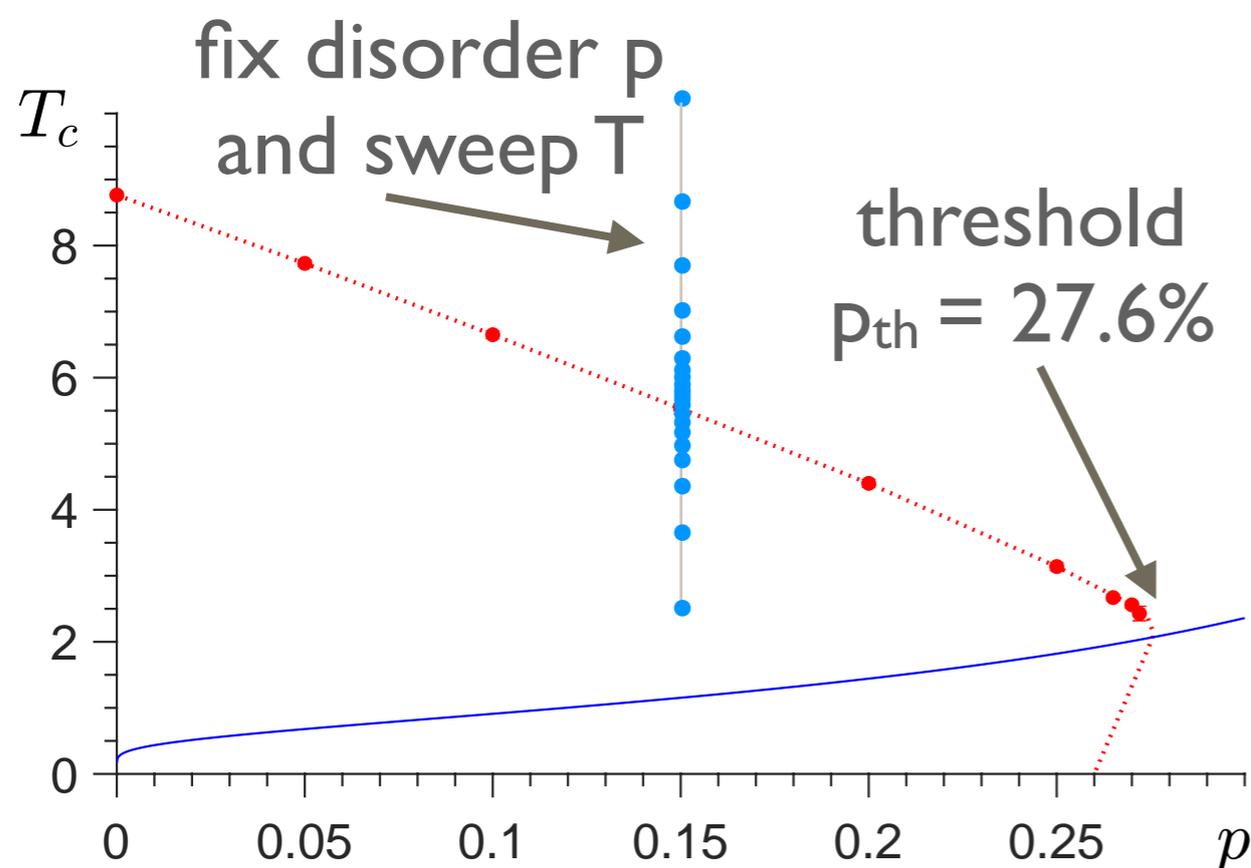


- For p=0 models are dual (low- and high-T expansions match).

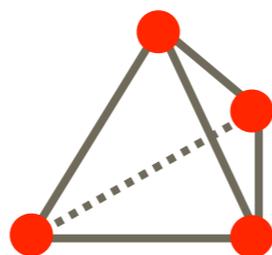
# NUMERICS



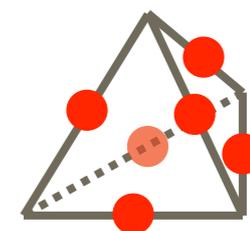
# 3D COLOR CODE THRESHOLDS FROM PHASE DIAGRAMS



- 4-body RCIM
- threshold for 2D sheet-like logical X



- 6-body RCIM
- threshold for 1D string-like logical Z



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# DISCUSSION

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- **Our results:**
  - local decoders of toric and color codes w/ provable thresholds,
  - noisy measurements — 3D TC sustainable threshold  $p_{TC}^{(2)} \approx 2\%$ .
  - 3D color code optimal thresholds from stat-mech:  
 $p^{(1)} \approx 1.9\%$  and  $p^{(2)} \approx 27.6\%$ .
- **3D gauge color code:** threshold  $p^{(1)}$  for 1D string-like (*Brown et al.'16*).

**THANK YOU FOR YOUR ATTENTION!**