

# Schur-Weyl Duality for the Clifford Group: Property testing, *de Finetti* representations, and a robust Hudson theorem

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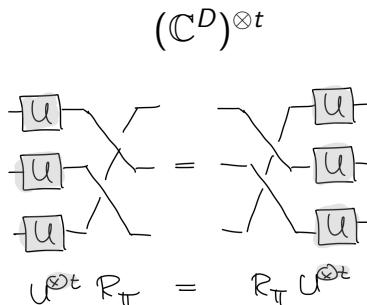


# Schur-Weyl duality

$$U^{\otimes t} |x_1, \dots, x_t\rangle = U |x_1\rangle \otimes \dots \otimes U |x_t\rangle$$

$$R_\pi |x_1, \dots, x_t\rangle = |x_{\pi^{-1}(1)}, \dots, x_{\pi^{-1}(t)}\rangle$$

**Schur-Weyl duality:** These actions generate each other's commutant.



Two *symmetries* that are ubiquitous in quantum information theory:

- ▶ **i.i.d. quantum information:**  $[\rho^{\otimes t}, R_\pi] = 0$
- ▶ eigenvalues, entropies,  $\dots$ :  $\rho \equiv U\rho U^\dagger$
- ▶ **randomized constructions:**  $E_{\text{Haar}}[|\psi\rangle\langle\psi|^{\otimes t}]$

See several other talks at this QIP...

## Clifford unitaries and stabilizer states

$$\mathbb{C}^D = (\mathbb{C}^d)^{\otimes n}$$

**Clifford group:** Unitaries  $U_C$  such that  $P \text{ Pauli} \Rightarrow U_C P U_C^\dagger \text{ Pauli}$ .  
For qubits, generated by CNOT, H, P.

**Stabilizer states:** States of the form  $|S\rangle = U_C |0\rangle^{\otimes n}$ .

These are important classes of unitaries & states:

- ▶ QEC, MBQC, topological order, randomized benchmarking, ...
- ▶ can be highly entangled, but efficient to represent and compute with
- ▶ 2-design; 3-design for qubits  $\Rightarrow$  efficient random constructions

*Motivates a Schur-Weyl duality for the Clifford group!*

# Our results

“Schur-Weyl duality” for the **Clifford group**: We characterize precisely which operators commute with  $U_C^{\otimes t}$  for all Clifford unitaries  $U_C$ .

*Fewer unitaries  $\leadsto$  larger commutant (more than permutations).*

Applications:

▶ **Property testing**

▶ **De Finetti theorems** with increased symmetry

▶ **Higher moments of stabilizer states**

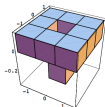
▶  **$t$ -designs** from Clifford orbits

▶ Robust **Hudson theorem**

$$|S\rangle^{\otimes t} \longleftrightarrow |\psi\rangle^{\otimes t}$$

$$\Psi_S \approx \sum_S p_S |S\rangle\langle S|^{\otimes S}$$

$$E_S[|S\rangle\langle S|^{\otimes t}]$$



# Towards Schur-Weyl duality for the Clifford group

Plan:

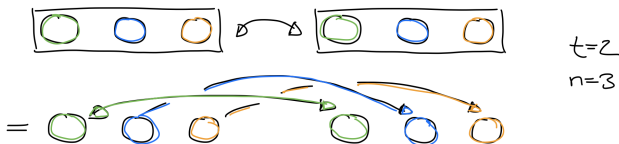
- 1 Write down permutation action in clever way.
- 2 Generalize.
- 3 Prove it!



# Towards Schur-Weyl duality for the Clifford group

- 1 Write down permutation action in clever way:

Permutation of  $t$  copies of  $(\mathbb{C}^d)^{\otimes n}$ :



$$R_{\pi} = r_{\pi}^{\otimes n}, \quad r_{\pi} = \sum_{\mathbf{x}} |\pi \mathbf{x}\rangle \langle \mathbf{x}|$$

Here, we think of  $\pi$  as  $t \times t$ -**permutation matrix**, and  $|\mathbf{x}\rangle = |x_1, \dots, x_t\rangle$  is computational basis of  $(\mathbb{C}^d)^{\otimes t}$ .

# Towards Schur-Weyl duality for the Clifford group

2 Generalize:

$$R_O = r_O^{\otimes n}, \quad r_O = \sum_{\mathbf{x}} |O\mathbf{x}\rangle \langle \mathbf{x}|$$

Allow all **orthogonal** and **stochastic**  $t \times t$ -matrices  $O$  with entries in  $\mathbb{F}_d$ .

For qubits, an example is the  $6 \times 6$  **anti-identity**:

$$\overline{\text{id}} = \begin{pmatrix} 0 & 1 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 0 \end{pmatrix},$$

$$R_{\overline{\text{id}}} |\mathbf{x}_1, \dots, \mathbf{x}_6\rangle = |\mathbf{x}_2 + \dots + \mathbf{x}_6, \dots, \mathbf{x}_1 + \dots + \mathbf{x}_5\rangle$$

The operator  $R_{\overline{\text{id}}}$  commutes with  $U_C^{\otimes 6}$  for every  $n$ -qubit Clifford unitary.

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$$R_{\overline{\text{id}}} |\mathbf{x}_1, \dots, \mathbf{x}_6\rangle = |\mathbf{x}_2 + \dots + \mathbf{x}_6, \dots, \mathbf{x}_1 + \dots + \mathbf{x}_5\rangle$$

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# Towards Schur-Weyl duality for the Clifford group

- 2 Generalize further:

$$R_T = r_T^{\otimes n}, \quad r_T = \sum_{(\mathbf{y}, \mathbf{x}) \in T} |\mathbf{y}\rangle \langle \mathbf{x}|$$

Allow all subspaces  $T \subseteq \mathbb{F}_d^{2t}$  that are **self-dual**, i.e.  $\mathbf{y} \cdot \mathbf{y}' = \mathbf{x} \cdot \mathbf{x}'^\dagger$  and of dimension  $t$ , and contain  $\mathbf{1} = (1, \dots, 1)$ .

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<sup>†</sup>For qubits, require modulo 4 ('doubly even' code).

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## Result

For  $n \geq t - 1$ , the operators  $R_T$  are  $\prod_{k=0}^{t-2} (d^k + 1)$  many linearly independent operators that span the commutant of  $\{U_{\mathbb{C}}^{\otimes t}\}$ .

*Sketch of proof:* Phase space formalism. Compute cardinalities and compare. □

*Independent of  $n!$  Rich algebraic structure (see paper).*

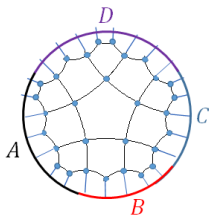


# Application 1: Higher moments of stabilizer states

Result ( $t$ -th moment)

$$E[|S\rangle\langle S|^{\otimes t}] \simeq \sum_{\mathcal{T}} R_{\mathcal{T}}$$

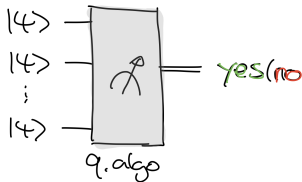
- ▶ When stabilizer states form  $t$ -design, reduces to  $\sum_{\pi} R_{\pi}$  (Haar average)
- ▶ Summarizes all previous results on statistical properties
- ▶ ... but also holds for larger  $t$ !



*We can also write  $t$ -th moment as weighted sum of certain CSS codes.*

## Application 2: Stabilizer testing

Given  $t$  copies of an unknown state in  $(\mathbb{C}^d)^{\otimes n}$ , decide if it is a stabilizer state or  $\epsilon$ -far from it.



*Idea:* Use the anti-identity. Measure POVM element  $\frac{1+R_{\text{id}}}{2}$  on  $t = 6$  copies.

### Result

Let  $\psi$  be a pure state of  $n$  qubits. If  $\psi$  is a stabilizer state then this accepts always. But if  $\max_S |\langle \psi | S \rangle|^2 \leq 1 - \epsilon^2$ , acceptance probability  $\leq 1 - \epsilon^2/4$ .

- ▶ Power of test independent of  $n$ . Answers q. by Montanaro & de Wolf.
- ▶ Similar result for qudits & for testing if blackbox unitary is Clifford.

*Why does it work? How to implement?*

# Stabilizer testing using Bell difference sampling

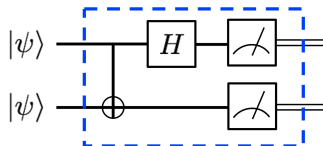
Any state  $\psi$  can be expanded in Pauli basis<sup>†</sup>:

$$\psi = \sum_{\mathbf{v}} c_{\psi} P_{\mathbf{v}}$$

- ▶ If **pure**, then  $p_{\psi}(\mathbf{v}) = 2^n |c_{\psi}(\mathbf{v})|^2$  is a probability distribution.
- ▶ If **stabilizer state**, then support of  $p_{\psi}$  is stabilizer group (up to sign).

**Key idea:** Sample & verify!

*How to sample?* If  $\psi$  is real, can simply measure in Bell basis ( $P_{\mathbf{v}} \otimes I$ )  $|\Phi^+\rangle$   
(**Bell sampling**; Montanaro, Zhao *et al*).

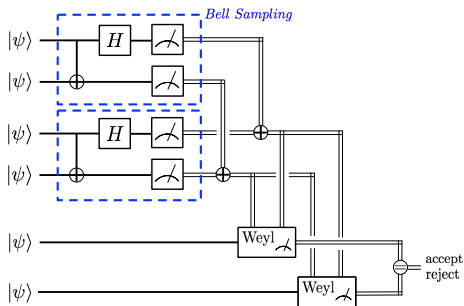


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<sup>†</sup> $P_{\mathbf{v}} = P_{v_1} \otimes \dots \otimes P_{v_n}$  where  $P_{00} = I$ ,  $P_{01} = X$ ,  $P_{10} = Z$ ,  $P_{11} = Y$

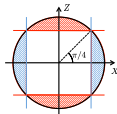
# Stabilizer testing using Bell difference sampling

In general, need to take **difference** of two Bell measurement outcomes:



- ▶ Constant-depth circuit, only need coherent two-qubit operations.
- ▶ Circuit is equivalent to measuring the anti-identity!

*Proof of converse uses **uncertainty relation**.*



## Application 3: Stabilizer de Finetti theorems

Any tensor power  $|\psi\rangle^{\otimes t}$  has  $S_t$ -symmetry. De Finetti theorems provide 'partial' converse: If  $|\Psi\rangle$  has  $S_t$ -symmetry,  $\Psi_s \approx \int d\mu(\psi)\psi^{\otimes s}$  for  $s \ll t$ .

Stabilizer tensor powers have **increased symmetry**:

$$R_O |S\rangle^{\otimes t} = |S\rangle^{\otimes t} \quad \text{for all orthogonal and stochastic } O$$

### Result

Assume that  $|\Psi\rangle \in ((\mathbb{C}^d)^{\otimes n})^{\otimes t}$  has this symmetry,  $d > 2$ . Then:

$$\|\Psi_s - \sum_S p_S |S\rangle\langle S|^{\otimes s}\|_1 \lesssim d^{2n(n+2)} d^{-(t-s)/2}$$

- ▶ Approximation is **exponentially good**, yet *bona fide* stabilizer states.
- ▶ Similar to Gaussian de Finetti (Leverrier *et al*). Applications to QKD?

*Can reduce symmetry requirements at expense of approximation.*

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## Application 4: $t$ -designs from Clifford orbits

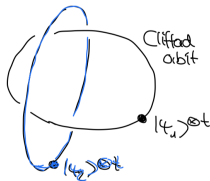
When  $t > 2$  or 3 (qubits), stabilizer states fail to be  $t$ -design. Yet, hints in the literature that this failure is relatively *graceful* (Zhu *et al*, Nezami-W). E.g., Clifford orbit of non-stabilizer qutrit states can be 3-design!

We prove in general:

### Result

For every  $t$ , there exists ensemble of  $N = N(d, t)$  many fiducial states in  $(\mathbb{C}^d)^{\otimes n}$  such that corresponding Clifford orbits form  $t$ -design.

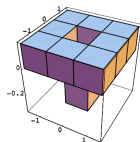
- ▶ Number of fiducials does not depend on  $n$ !
- ▶ Efficient construction?



## Application 5: Robust Hudson theorem

For odd  $d$ , every quantum state has a discrete **Wigner function**:

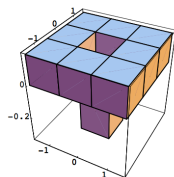
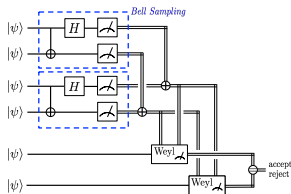
$$W_\rho(\mathbf{v}) = d^{-2n} \sum_{\mathbf{w}} e^{-2\pi i[\mathbf{v}, \mathbf{w}]/d} \text{tr}[\rho P_{\mathbf{v}}]$$



- ▶ Quasi-probability distribution on phase space  $\mathbb{F}_d^{2n}$
- ▶ **Discrete Hudson theorem**: For pure states,  $W_\psi \geq 0$  iff  $\psi$  stabilizer
- ▶ Wigner negativity  $\text{sn}(\psi) = \sum_{\mathbf{v}: W_\rho(\mathbf{v}) < 0} |W_\rho(\mathbf{v})|$ : monotone in resource theory of stabilizer computation; witness for contextuality

### Result (Robust Hudson)

There exists a stabilizer state  $|S\rangle$  such that  $|\langle S|\psi\rangle|^2 \geq 1 - 9d^2 \text{sn}(\psi)$ .



Schur-Weyl duality for the Clifford group:

- ▶ clean algebraic description in terms of self-dual codes
- ▶ resolve open question in quantum property testing
- ▶ new tools for stabilizer states: moments, de Finetti, Hudson, ...

*Thank you for your attention!*