

# Non-closure of the set of quantum correlations via graphs

Jitendra Prakash

Department of Pure Mathematics and Institute for Quantum Computing  
University of Waterloo

(Joint work with Vern I. Paulsen and Ken Dykema)

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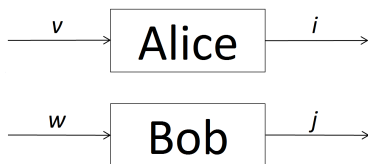
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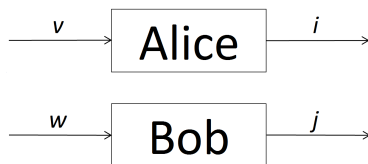
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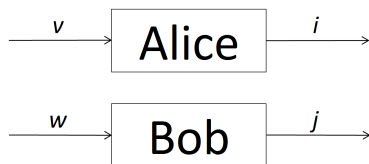
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Tsirelson [Ts1993] considered different mathematical models to describe these correlations and studied relationships among them.

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$$\rho(0, 0|v, v) = t, \quad \rho(0, 1|v, v) = \rho(1, 0|v, v) = 0, \quad \rho(1, 1|v, v) = 1 - t,$$

$$\rho(0, 0|v, w) = \frac{t(5t-1)}{4}, \quad \rho(1, 1|v, w) = \frac{(1-t)(4-5t)}{4},$$

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Then  $(\rho(i, j|v, w)) \in \overline{C_q(5, 2)} \setminus C_q(5, 2)$

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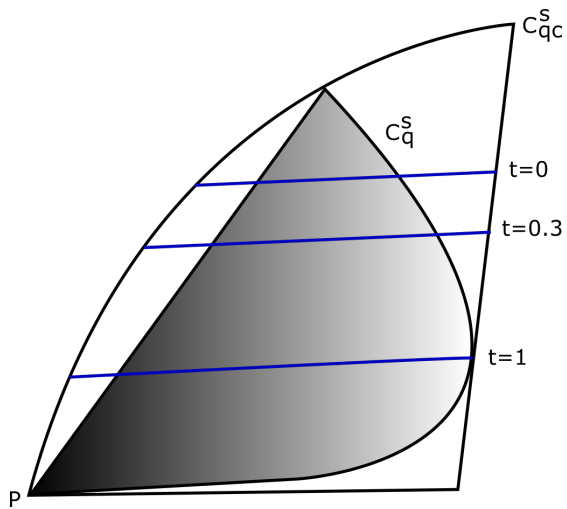
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## Theorem (PSSTW)

A correlation  $(p(i, j|v, w))$  belongs to  $C_{qc}^s(n, k)$  (resp.  $C_q^s(n, k)$ ) if and only if there exists a unital  $C^*$ -algebra  $\mathcal{A}$  (resp. finite dimensional  $C^*$ -algebra) with a tracial state  $\tau$  and projections  $\{e_{v,i} : 1 \leq v \leq n, 1 \leq i \leq k\}$  such that  $\sum_i e_{v,i} = 1$  and

$$p(i, j|v, w) = \tau(e_{v,i} e_{w,j})$$



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By the above characterisation,

$$f_q(t) = \inf \left\{ \sum_{(v,w) \in E} \tau(e_v e_w) : e_v \in \mathcal{A} \text{ projections, } \tau(e_v) = t \forall v \in V, \right.$$

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*Let  $G = (V, E)$  be a vertex and edge transitive graph on  $n$  vertices and let  $t \in [0, 1]$  be irrational. If  $f_q(t)$  attains the infimum then there exists a nondegenerate interval  $[r, s]$  having rational endpoints such that  $t \in [r, s]$  and the restriction of  $f_q$  to  $[r, s]$  is linear.*

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### Corollary ( $G = K_5$ )

*If  $C_q(5, 2)$  is closed, then for each irrational  $t \in [0, 1]$  there exists a nondegenerate interval  $[r, s]$  with  $t \in [r, s]$  such that  $f_q|_{[r, s]}$  is linear.*

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### Theorem (Kruglyak, Rabanovich, Samaoilenko)

Let  $n \geq 5$  and let  $\alpha \in (\beta_n, n - \beta_n)$ . Then there exist finite dimensional projections  $P_1, \dots, P_n$  such that  $\sum_{j=1}^n P_j = \alpha I$  if and only if  $\alpha \in \mathbb{Q}$ .

$$G = K_5$$

Theorem (Dykema, Paulsen, P.)

$C_q^s(5, 2)$  is not closed.

- 1 Show that

$$f_q(t) \geq 5t(5t - 1) \text{ if } t \in \left[ \frac{1}{5}, \frac{4}{5} \right].$$

- 2 Show that  $f_q(t) = 5t(5t - 1)$  for all rational  $t \in \frac{1}{5}[\beta_5, 5 - \beta_5] \subset \left[ \frac{1}{5}, \frac{4}{5} \right]$ . Enough to show that  $f_q(t) \leq 5t(5t - 1)$  on that interval.

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$$\text{tr}_{5k}(\tilde{P}_i) = \frac{1}{5k} \sum_{j=1}^5 \text{Tr}(P_j) = \frac{1}{5k} \text{Tr} \left( \sum_{j=1}^5 P_j \right) = \frac{1}{5k} (5tk) = t.$$

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- 3 Contradicts “piecewise” linearity.

Thank You.