



# Quantum fully homomorphic encryption with verification



JOINT CENTER FOR  
QUANTUM INFORMATION  
AND COMPUTER SCIENCE



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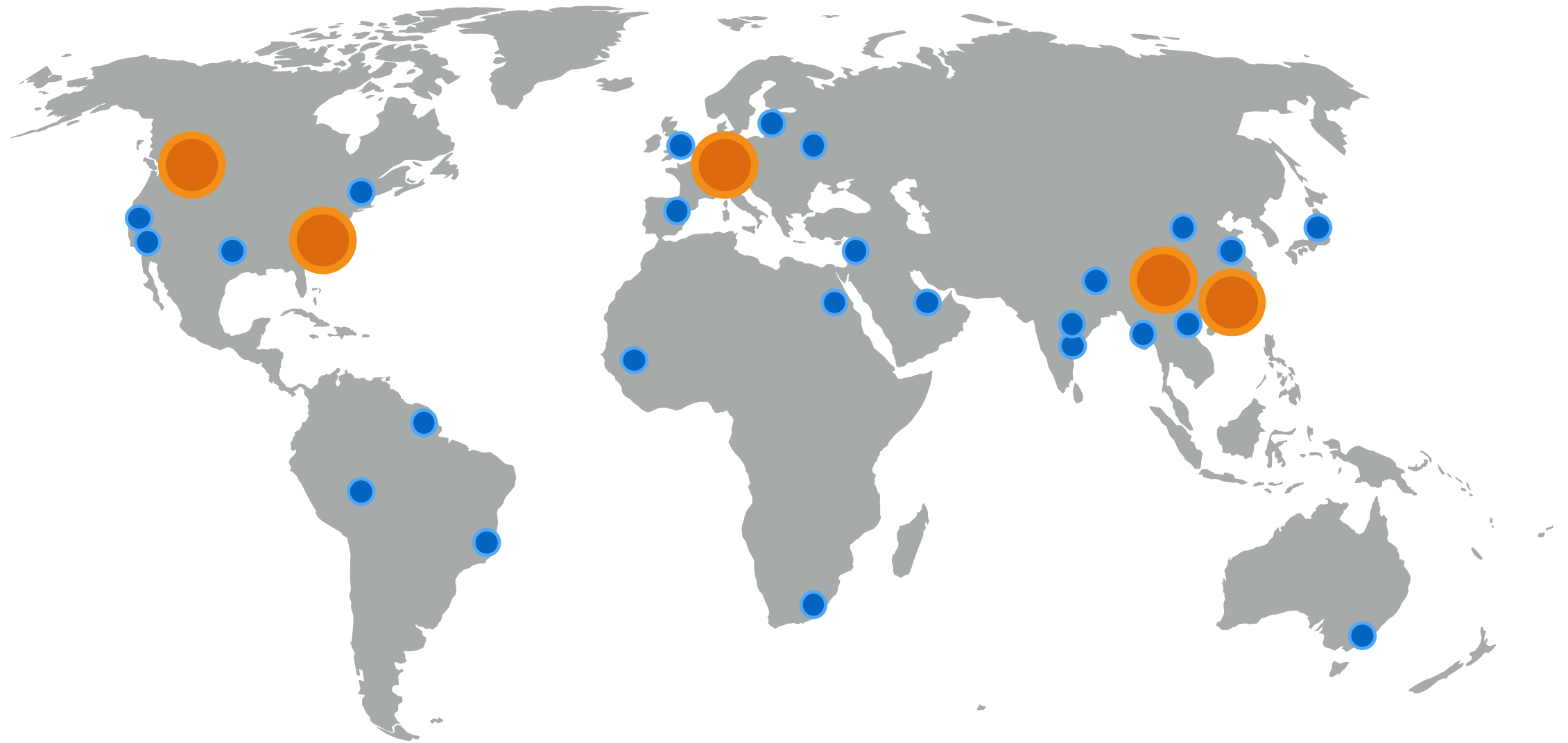
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and Florian Speelman

arXiv:1708.09156

# A Possible Future for Quantum Computing







private — efficient — verifiable

Disclaimer: dots placed carelessly

# Outline

1. (Quantum) Fully Homomorphic Encryption
2. Classical application: zero-knowledge proofs
3. Our contributions:
  - Definitions
  - Construction
  - Application

# Classical Fully Homomorphic Encryption

	KeyGen	$k, evk$
	Enc	$k, x \rightarrow \bar{x}$
	Eval	$evk, f, \bar{x} \rightarrow \bar{y}$
	Dec	$k, \bar{y} \rightarrow y$

**Homomorphic** if  $y = f(x)$

**Classical FHE** possible [Gen09] under comp assumptions.  
Current focus: efficiency

# Quantum Fully Homomorphic Encryption



KeyGen

$k, evk$



Enc

$k, \sigma \rightarrow \bar{\sigma}$



Eval

$evk, C, \bar{\sigma} \rightarrow \bar{\rho}$



Dec

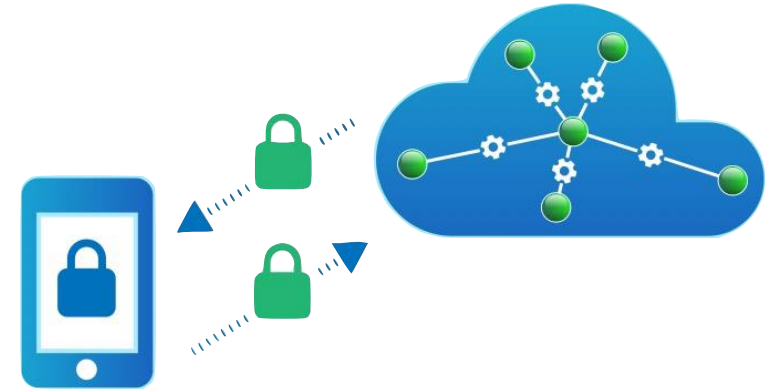
$k, \bar{\rho} \rightarrow \rho$

**Homomorphic** if  $\rho = C(\sigma)$

Quantum FHE possible [DSS16] under comp assumptions.  
Current focus: efficiency

# Applications of Fully Homomorphic Encryption

**Practical application:**  
outsourcing computations



**Theoretical applications (classical):**  
multiparty computation, functional encryption,  
private information retrieval, zero-knowledge proofs,  
and more...

**Theoretical applications (quantum): ???**

# Zero-knowledge proofs (classical)

SAT: given a formula  $\phi$ , is  $\phi$  satisfiable?

I know!  $\phi$  is satisfiable!

Wonderful! Can you prove it to me?  
I am only a polynomial-time human...

Well, I know a satisfying assignment  $w$ ,  
but I am not going to tell YOU...

Let's do a zero-knowledge proof!





# Zero-knowledge from FHE (classical)

$$\bar{w} \leftarrow \text{Enc}_k(w)$$

$$b \in_R \{0,1\}$$

if  $b = 1$ :  $\bar{c} \leftarrow \text{Eval}_{evk}(\text{witness-check-fn}, \bar{w})$   
if  $b = 0$ :  $\bar{c} \leftarrow \text{Eval}_{evk}(\text{set-to-0}, \bar{w})$

$$c \leftarrow \text{Dec}_k(\bar{c})$$

if  $c = b$ , **accept** (otherwise **reject**)

(*prover*)



(*verifier*)

Prover cannot cheat... but verifier can!



# Zero-knowledge from FHE (classical)

$$\bar{w} \leftarrow \text{Enc}_k(w)$$

$$b \in_R \{0,1\}$$

if  $b = 1$ :  $\bar{c} \leftarrow \text{Eval}_{evk}(\text{witness-check-fn}, \bar{w})$   
if  $b = 0$ :  $\bar{c} \leftarrow \text{Eval}_{evk}(\text{set-to-0}, \bar{w})$

$$\text{commit to } c \leftarrow \text{Dec}_k(\bar{c})$$

proof of computation (transcript)

if transcript is ok, **reveal**  $c$

if  $c = b$ , **accept** (otherwise **reject**)

(*prover*)



(*verifier*)

# Verification in (Q)FHE

- ▶ **Verification of computation** is crucial in applications
- ▶ **Classical FHE** has verification “automatically”
- ▶ **Quantum FHE** does not

# Verification in (Q)FHE

- ▼ **Verification of computation** is crucial in applications



I will only accept your output if I can verify that you applied the right circuit to my input!



- ▶ **Classical FHE** has verification “automatically”
- ▶ **Quantum FHE** does not

# Verification in (Q)FHE

- ▶ **Verification of computation** is crucial in applications
- ▼ **Classical FHE** has verification “automatically”



Here is a transcript of all the steps in my computation, please check them all.



- ▶ **Quantum FHE** does not

# Verification in (Q)FHE

- ▶ **Verification of computation** is crucial in applications
- ▶ **Classical FHE** has verification “automatically”
- ▼ **Quantum FHE** does not  
Measurement and no-cloning  
prevent this easy solution



The outcome of that measurement was 0, truly! You must believe me!



# Related topics

- ▶ **Quantum authentication**
- ▶ **Quantum computing on authenticated data**
- ▶ **Quantum fully homomorphic encryption**

# Related topics

- ▼ **Quantum authentication**
  - “verifiable HE” for the identity circuit
  - e.g. polynomial code [BCG+06], Clifford code [ABE10], trap code [BGS12]



Hold this qubit, I will be *right* back!



- ▶ **Quantum computing on authenticated data**
- ▶ **Quantum fully homomorphic encryption**



# Related topics

- ▶ **Quantum authentication**
- ▼ **Quantum computing on authenticated data**
  - interaction during evaluation
  - verification



- ▶ **Quantum fully homomorphic encryption**

# Related topics

- ▶ Quantum authentication
- ▶ Quantum computing on authenticated data
- ▼ Quantum fully homomorphic encryption
  - no interaction during evaluation
  - no verification



The result of your computation is 42.

Ok!



# Our contributions

**Definitions**

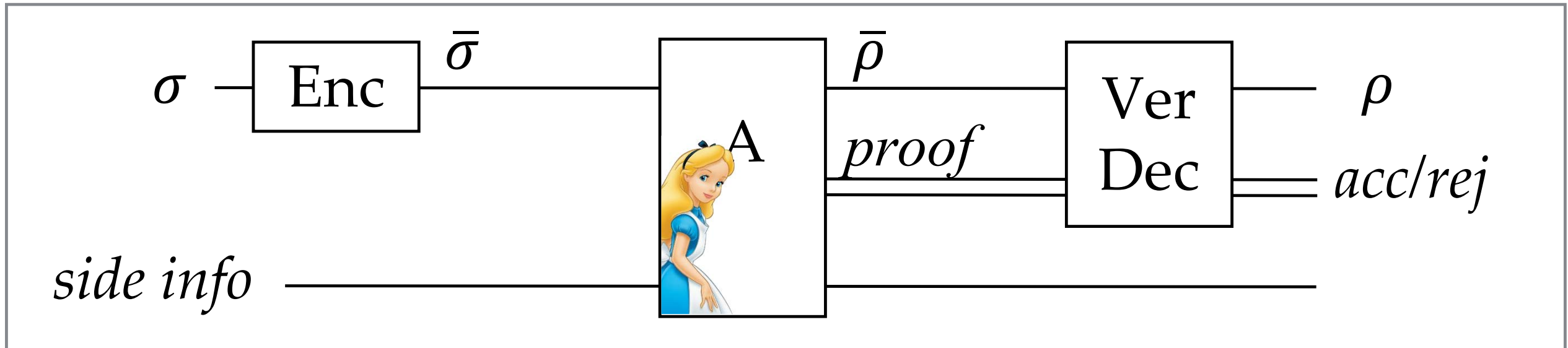
**Construction: a verifiable QFHE scheme**

**Application: quantum one-time programs**

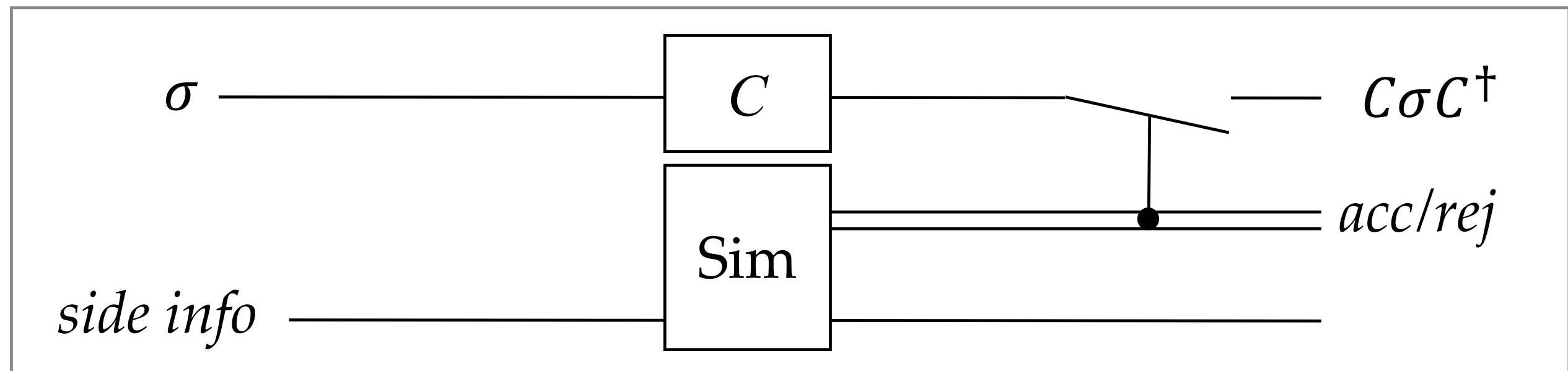
# Defns: Verification in QFHE

Dec replaced by **VerDec**:  $k, \bar{\rho}, C, \textit{proof} \rightarrow \rho, \textit{acc/rej}$

REAL



IDEAL



**Verifiability:**  $\forall \text{poly } A \exists \text{poly } \textit{Sim} \forall \sigma, \textit{side info}, C:$   
 REAL and IDEAL are indistinguishable

# Defns: Verification in QFHE

## Variant (indistinguishability)

- Alternative security definition in terms of a guessing game
  - Adversary has to guess whether he interacts with actual or idealized functionality
- Proven equivalent to semantic definition

## Compactness

- Verification: (classical) proof, output accept/reject
- Decryption: quantum input, runtime should **not** depend on circuit

# Construction: Verifiable QFHE

## Ingredients

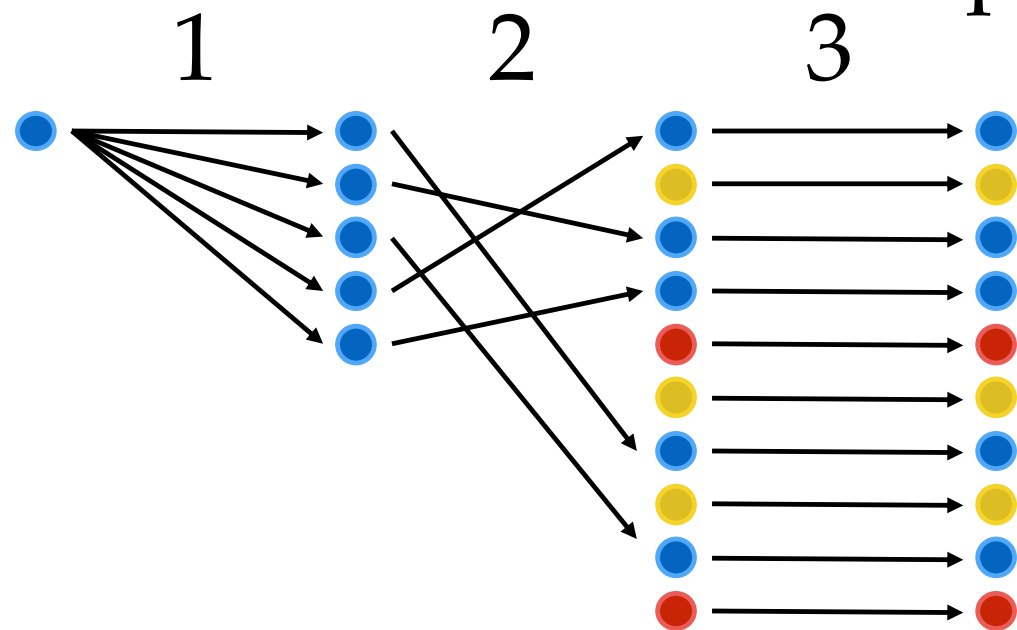
- Quantum authentication code: trap code
- Classical FHE: any quantum-secure scheme with low-depth decryption
- Classical MAC: any quantum-secure scheme

# Construction: Verifiable QFHE

traps:  $|0\rangle, |+\rangle$

## Encryption

Authenticate with trap code



1. error-correcting code
  2. traps in random positions
  3. quantum one-time pad
- secret key:** positions + pad keys

Encode secret key with classical FHE+MAC

Encrypted qubit:

(classical encrypted info, trap code qubits)



# Construction: Verifiable QFHE

## Evaluation

Clifford operations ( $X, Z, H, P, \text{CNOT}$ ):

Apply transversally (magic states)

Update encoded keys using FHE

T gate:

After applying, unwanted  $P$  error

Use (new, extended) [gadget](#) to remove the error

(using ideas from [DSS16])

## Verify & Decrypt

Check (classical) MAC+FHE transcript of key updates

Decrypt trap code if everything checks out

# Application: one-time programs

- Idea: Programs that ‘self-destruct’  
after a single execution
- Ingredients:
  - Classical one-time program
  - Verifiable QFHE scheme
- Simple construction
  - Q-OTP for  $C$ :  $(evk, \text{OTP for } \text{VerDec}(C, k, \cdot), \text{Enc}_k(C))$
- Not a new result [BGS12]



Proof of correct execution is classical



# Summary

## ▼ Definitions

- verifiability (semantic)
- verifiability (indistinguishability)
- compactness

▶ **Construction: a verifiable QFHE scheme**

▶ **Application: quantum one-time programs**

# Summary

- ▶ **Definitions**
- ▼ **Construction: a verifiable QFHE scheme**

Ingredients:

- classical FHE
- classical authentication code (MAC)
- quantum authentication code (trap code)

- ▶ **Application: quantum one-time programs**

# Summary

- ▶ **Definitions**
- ▶ **Construction: a verifiable QFHE scheme**
- ▼ **Application: quantum one-time programs**  
Alternative construction to [BGS12]. Ingredients:
  - classical one-time programs
  - QFHE with verification

# Future work / open questions

- Apply verifiable QFHE to build other advanced cryptographic primitives
- Are non-leveled schemes possible?
- Reduce client quantum capabilities

Thank you