

Single-shot security for one-time memories in the isolated qubits model

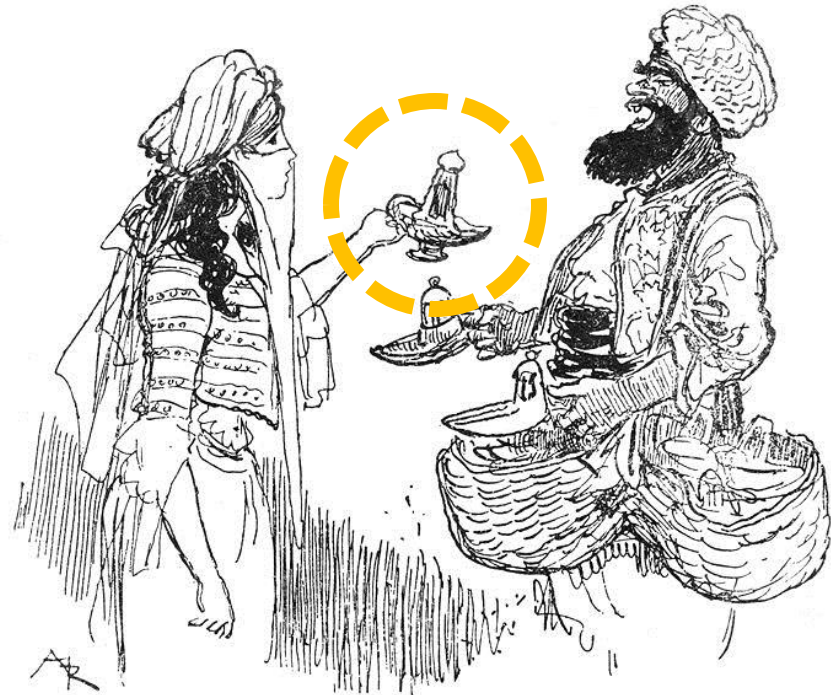
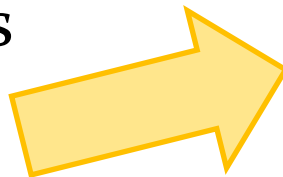
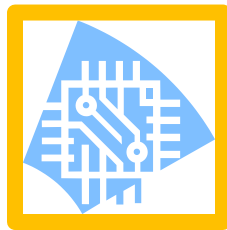
Yi-Kai Liu

National Institute of Standards and Technology (NIST)
Gaithersburg, MD, USA

Arxiv: 1402.0049

One-time memories

- Tamper-resistant cryptographic hardware
 - Needed in situations where Alice's data resides on hardware that is controlled by Eve
 - E.g., a stolen smartphone
- Want to use *simple* tamper-resistant chips to implement complex functions



One-time memories

- **One-time memory** (OTM) contains two messages s, t
 - Adversary can choose to read s or t , but not both
 - “Non-interactive oblivious transfer”
- Can be used to construct **one-time programs**
 - Evaluate some circuit
 - Can only be run *once*
 - Intermediate results of computation are hidden
 - [Goldwasser, Kalai & Rothblum, 2008], [Goyal et al, 2010]

One-time memories

- Can we build OTM's based on some physical principle?
 - Classical physics: no!
(information can always be copied)
 - Quantum physics: no!
(no-go theorems for bit-commitment, oblivious transfer)
- However, if one assumes that the adversary is **k-local**, then quantum bit-commitment is possible! [Salvail '98]
 - Adversary cannot entangle more than k qubits

Isolated qubits model

- All parties (both honest and dishonest) are restricted to **LOCC operations**
 - LOCC = “local operations and classical communication”
 - Pick a qubit, measure it, get some classical outcome, repeat...
 - No entangling gates
- Example: nuclear spins?
 - Isolated qubits can exist in a world with quantum computers



Isolated qubits model

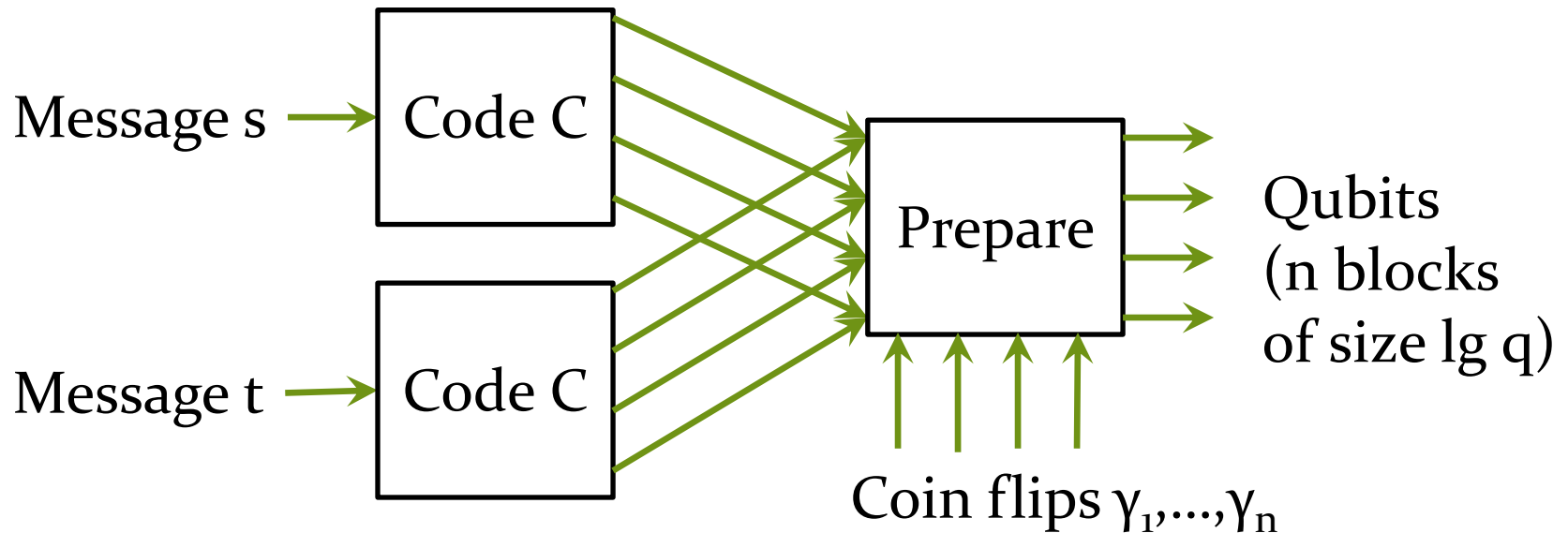
- Is there anything quantum going on here?
 - State remains separable at all times
- “Nonlocality without entanglement” [Bennett et al, 1999]
 - Certain transformations using LOCC operations *can be inverted* using entangling operations, *but cannot be inverted* using LOCC

Our results

New paper:
Arxiv:1402.0049

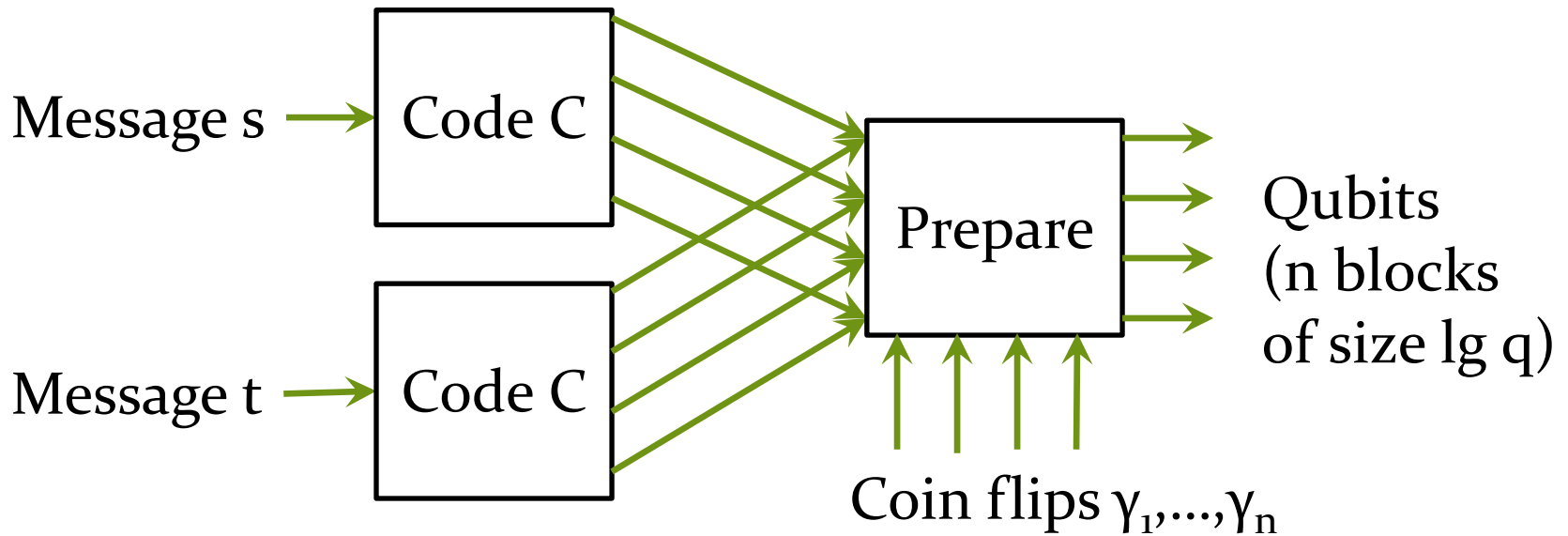
- One-time memories in the isolated qubits model
 - Based on Wiesner's idea of conjugate coding
 - **Single-shot security:** measure the adversary's uncertainty using the smoothed min-entropy
 - **Secure against general LOCC adversaries:** including adaptive sequences of weak measurements
 - **Efficiently implementable:** OTM's can be built using a large class of error-correcting codes

Conjugate coding



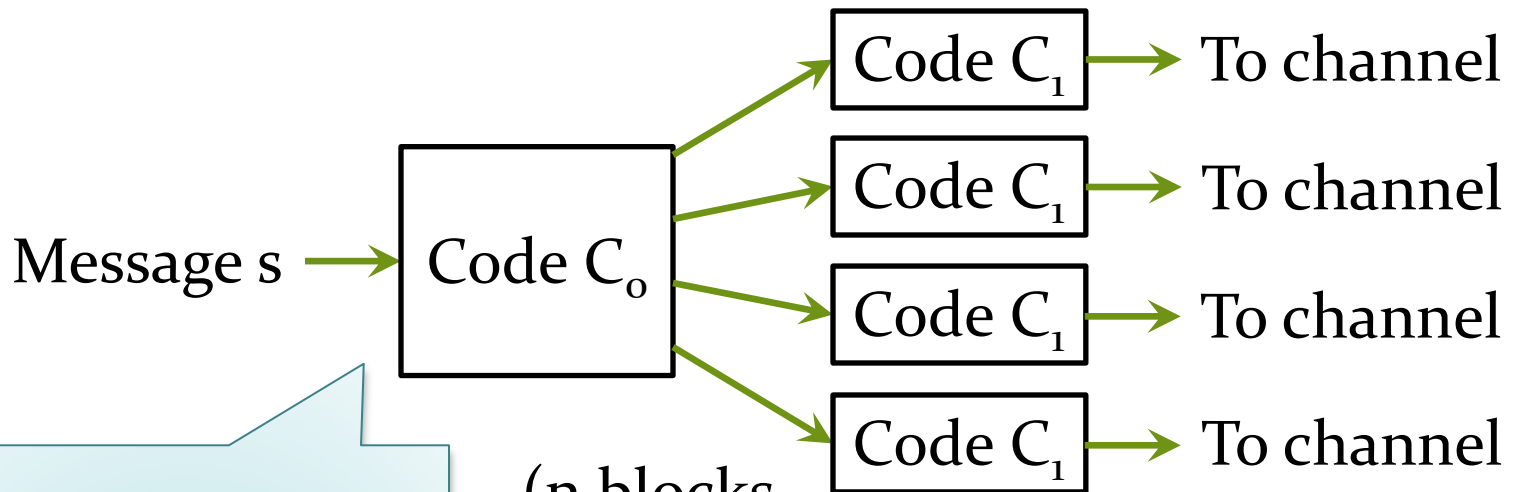
- To prepare the i 'th block of qubits:
- If $\gamma_i = 0$, use the i 'th block of $C(s)$ and the **standard basis**
- If $\gamma_i = 1$, use the i 'th block of $C(t)$ and the **Hadamard basis**

Conjugate coding



- To read **s**: measure qubits in **standard basis**
- To read **t**: measure qubits in **Hadamard basis**
- This is equivalent to receiving $C(s)$ or $C(t)$ through a **q-ary symmetric channel**

Good codes for the q-ary symmetric channel



(n blocks of size $\lg q_0$)

(n blocks of size $\lg q$)

Random binary linear code
Corrects erasure errors

Fixed binary linear code
Detects q-ary symmetric errors

Good codes for the q -ary symmetric channel

- For large q (growing with n), this approaches the capacity of the q -ary symmetric channel
- Efficient decoding: solving linear systems of equations over $GF(2)$
- Other constructions:
interleaved Reed-Solomon codes, interleaved AG codes
[Bleichenbacher et al; Shokrollahi; Brown et al]

Security

- **Ideal security goal:** adversary can learn either S or T, but not both
 - Impossible, if the adversary can perform entangling gates
- **We show a weaker (“leaky”) notion of security, in the isolated qubits model**
 - “Any cheating strategy *requires* entangling gates”
 - Honest strategies require only LOCC operations
 - However, some extra information leaks out
- For any LOCC adversary, $H_{\infty}^{\epsilon}(S, T | Z) \geq (0.5 - \delta) \ell$
 - Each of the messages S and T is ℓ bits long
 - Z is the adversary’s output

Security

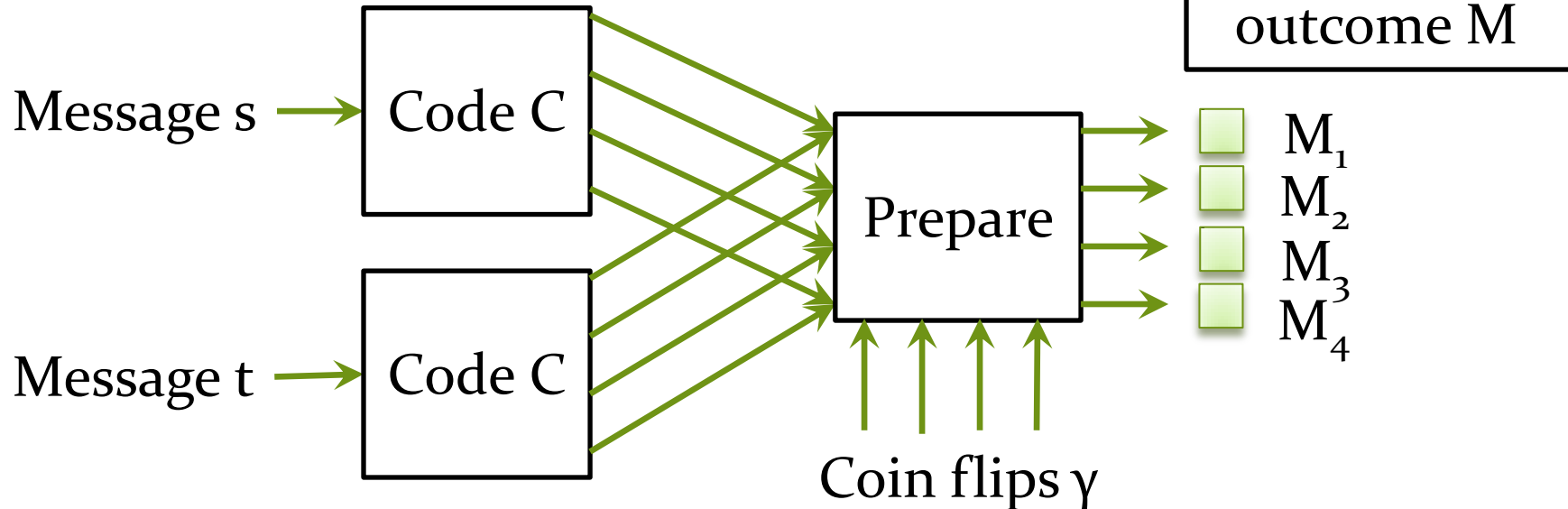
- Some issues to consider:
- **Privacy amplification** doesn't work in this setting
 - Honest parties can try to use a randomness extractor, but adversary also knows the seed!
- Security comes from the **choice of the code C**
 - Want it to be “unstructured” – what does this mean?
- General LOCC adversaries can be quite **complicated**
 - Can make a long sequence of weak measurements, w/ adaptive choices

Security proof

- Prove security against **separable** adversaries
 - Every POVM element is a tensor product of 1-qubit operators
 - Includes LOCC as a special case
- Assume the code C is **linear over $GF(2)$**
 - Given a random codeword, a large subset of the bits will be uniformly distributed => “unstructured”
 - Prevents the adversary from learning the basis choices γ
- Use a high-order **entropic uncertainty relation**
 - Measuring an arbitrary state in a random BB84 basis
 - Borrowed from the bounded quantum storage model [Damgard et al, 2006]

Given a **separable adversary A**

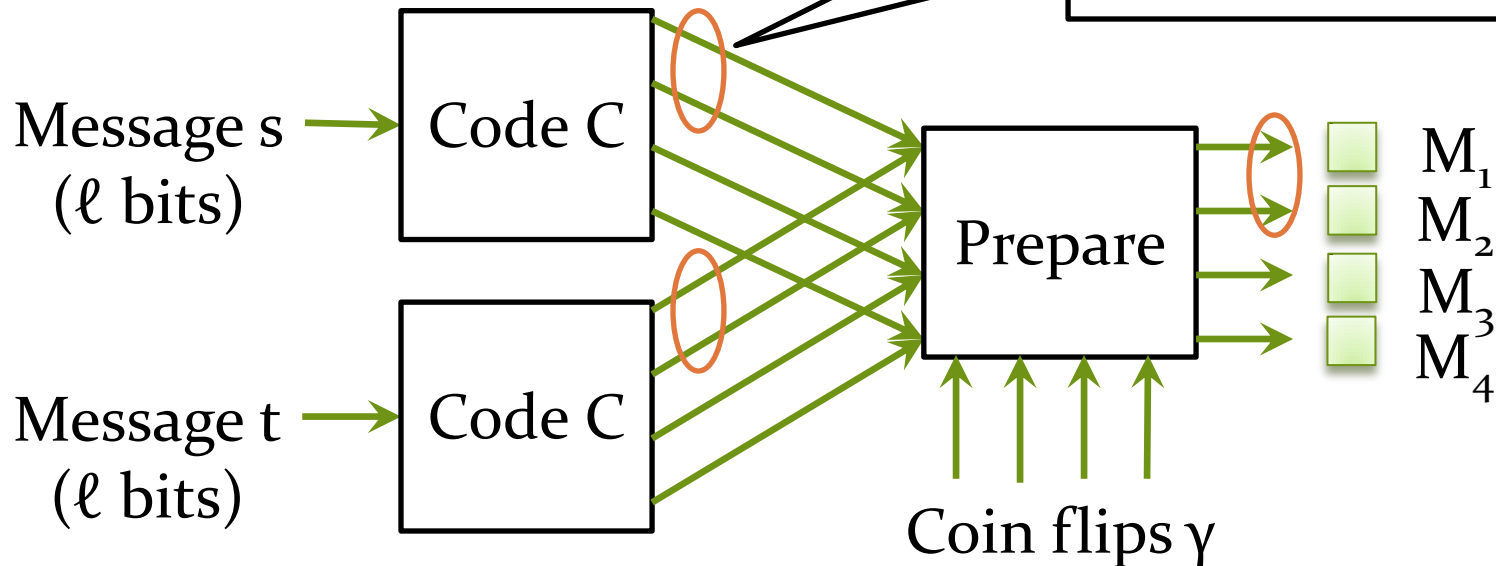
Security proof



- Want to analyze $\Pr(S, T | M)$
- Consider a **fictitious adversary A'** that measures each qubit once, and observes M_1, M_2, M_3, \dots (call this event M')
- Then $\Pr(S, T | M) = \Pr(S, T | M')$

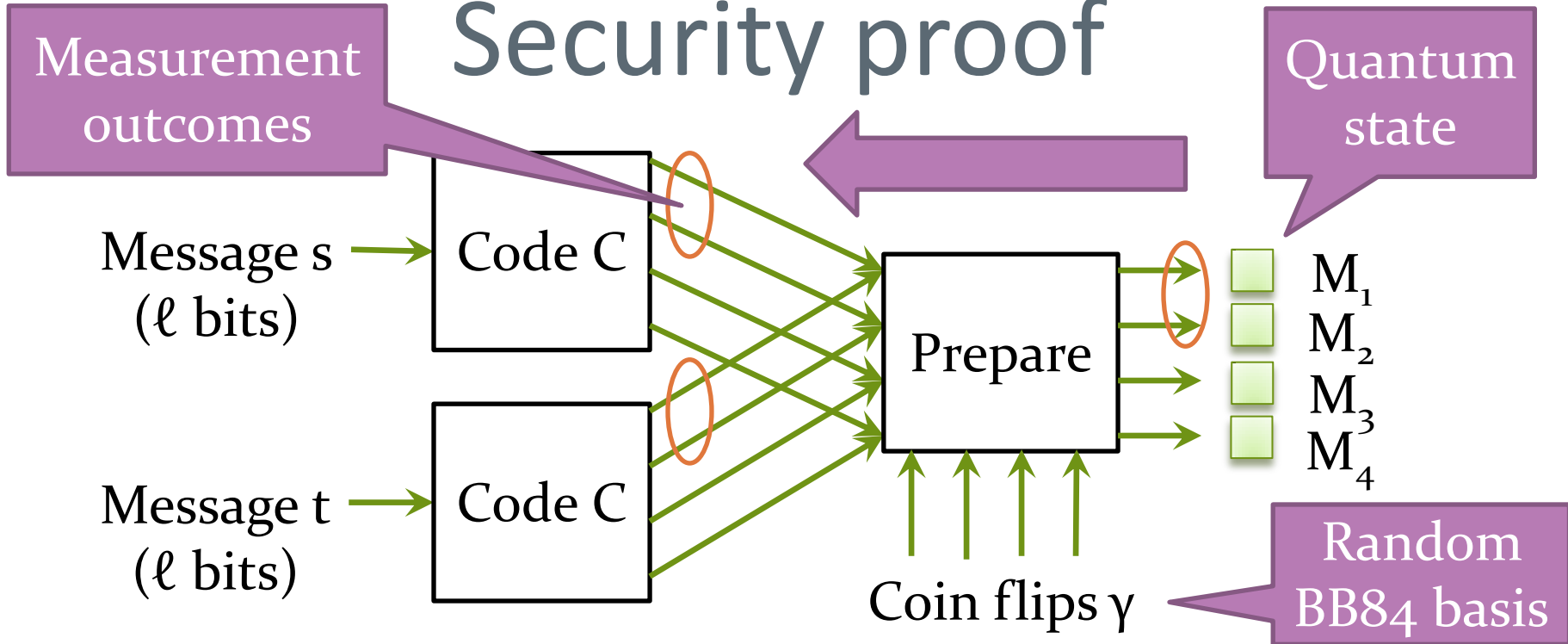
Security pro

There exists a subset of ℓ bits of $C(s)$ that are uniformly distributed



- Wlog, suppose the fictitious adversary A' **measures this subset of qubits first**, and observes M_1, M_2 (call this event M'')
- Want to analyze $\Pr(S, T | M'')$
- Note: coin flips Γ conditioned on M'' are still **uniformly distributed**

Security proof



- Note: coin flips Γ conditioned on M'' are still **uniformly distributed**
- Now run the experiment backwards...
- Use the uncertainty relation to lower-bound $H_\infty^\epsilon(S, T | M'')$

Outlook

- Isolated qubits model
- One-time memories (OTM's) using conjugate coding
 - Efficient implementations
 - Single-shot security against general LOCC adversaries
- **Can we control the leakage of information from our OTM's?**
 - Necessary to construct **one-time programs**
 - Note: LOCC also implies strong constraints on the **types of information** that the adversary can learn
 - Conjecture: for one-time programs based on **garbled circuits**, the relevant information cannot be extracted via LOCC
 - More generally, **can we construct ideal OTM's** using a random oracle, or some variant of leakage-resilient encryption?
- **Beyond LOCC and the isolated qubits model**
 - Are our OTM's secure against Salvail's k -local adversaries?