

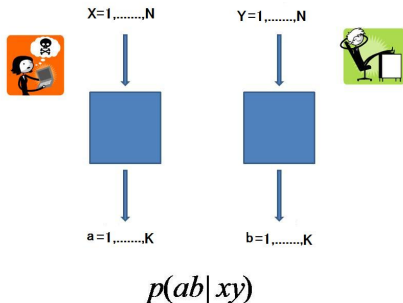
# Superactivation of quantum nonlocality

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# What is quantum nonlocality?



**Figure:** Alice and Bob measurements. Inputs:  $x$  and  $y$ , Outputs:  $a$  and  $b$ .  $P(a, b|x, y)$  is the probability of obtaining the pair  $(a, b)$  when Alice and Bob measure, respectively, with the input  $x$  and  $y$ .

We deal with the “**probability distributions**”

$$P = (P(a, b|x, y))_{x, y=1, \dots, N}^{a, b=1, \dots, K}$$

## What is quantum nonlocality?

- $P = \{P(a, b|x, y)\}_{x,y;a,b}$  is a **Classical prob. distribution** ( $P \in \mathcal{L}$ ) if it is in the convex hull of the elements of the form

$$P(a, b|x, y) = P_1(a|x)P_2(b|y) \text{ for every } x, y, a, b, \text{ where}$$

$$P_1(a|x) \geq 0 \text{ and } \sum_a P_1(a|x) = 1 \text{ for every } x \text{ (similar for } P_2).$$

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- $P$  is a **Quantum prob. distribution** ( $P \in \mathcal{Q}$ ) if it can be written like:

$$P(a, b|x, y) = \text{tr}(E_x^a \otimes F_y^b \rho) \text{ for every } x, y, a, b, \text{ where}$$

$\rho$  is a bipartite quantum state and  $\{E_x^a\}_a$  is a POVM for every  $x$  (similar for  $\{F_y^b\}_b$ ).

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- **Quantum nonlocality:**  $\mathcal{L} \subsetneq \mathcal{Q}$ .

# Why is quantum nonlocality interesting?

- Fundamental and fascinating phenomenon!

IDEA: THEORY  $\leftrightarrow$  EXPERIMENTAL DATA

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- It has many applications:
  1. Device Independent Quantum Cryptography.
  2. Generation of random numbers.
  3. Interactive proof systems.

... ..

## How to study nonlocality? Bell inequalities

- Given  $M = \{M_{x,y}^{a,b}\}_{x,y;a,b}$  and  $P = \{P(a,b|x,y)\}_{x,y;a,b}$ , denote

$$\langle M, P \rangle = \sum_{x,y,a,b} M_{x,y}^{a,b} P(a,b|x,y) \quad \text{and} \quad \omega(M) = \sup_{P \in \mathcal{L}} |\langle M, P \rangle|.$$



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- How nonlocal is  $\rho$ ? Define

$$\mathcal{Q}_\rho := \{P = \{tr(E_x^a \otimes F_y^b \rho)\} : \{E_x^a\}, \{F_y^b\} \text{ POVMs.}\}$$

Then,

$$LV(\rho) = \sup_{P \in \mathcal{Q}_\rho} \nu(P).$$

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- $LV(\rho) \geq 1$  for ever  $\rho$ . Moreover,  $\rho$  local  $\Leftrightarrow LV(\rho) = 1$ .

We are interested in:

- Behavior of  $LV(\rho)$ ... asymptotically .... at least for some states...

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- Superactivation of nonlocality:

Find a local state  $\rho$  such that  $LV(\rho^{\otimes k}) > 1$  for some  $k$ .

## An upper bound for $LV(\rho)$

- Theorem: Given any pure state  $|\psi\rangle = \sum_{i=1}^n \alpha_i |ii\rangle$ ,  $\alpha_i \geq 0$  we have

$$LV(|\psi\rangle) \leq UB(|\psi\rangle) := \left( \sum_{i=1}^n \alpha_i \right)^2.$$

In particular,  $LV(\rho) \leq n$  for every  $n$  dimensional quantum state  $\rho$ .

Note that  $UB(|\varphi_n\rangle) = n$ .

## Lower bound and the Khot and Visnoi game

- In a previous work, Buhrman, Regev, Scarpa and de Wolf showed that for every  $n$ ,

$$LV(|\varphi_n\rangle) \geq C \frac{n}{(\log n)^2}.$$

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- There exist a Bell inequality  $M(n) = \{M(n)_{x,y}^{a,b}\}_{x,y,a,b}$  and some POVMs  $\{E(n)_x^a\}_{x,a}$ ,  $\{F(n)_y^b\}_{y,b}$  such that

$$\frac{1}{\omega(M(n))} \sum_{x,y,a,b} M(n)_{x,y}^{a,b} \langle \varphi_n | E(n)_x^a \otimes F(n)_y^b | \varphi_n \rangle \geq C \frac{n}{(\log n)^2}.$$

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- In fact,

$$C \frac{UB(|\psi\rangle)}{(\log n)^2} \leq LV(|\psi\rangle) \leq UB(|\psi\rangle)$$

for every  $n$  dimensional pure state  $|\psi\rangle$  (even better!).

## A tighter upper bound

- Theorem: For every  $n$ ,

$$LV(|\varphi_n\rangle) \leq D \frac{n}{\sqrt{\log n}}.$$

This shows

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- We understand better the quantity  $LV(|\varphi_n\rangle)$ .
- Interesting part of the result: Proof

How is this related with multiplicativity?

- An easy proof of  $\frac{LV(|\varphi_n\rangle^{\otimes 5})}{LV(|\varphi_n\rangle)^5} \geq C' \sqrt{\log n}$ :

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- On the other hand, applying the KV game we have

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- Therefore,

$$\frac{LV(|\varphi_n\rangle^{\otimes 5})}{LV(|\varphi_n\rangle)^5} \geq \frac{C n^5 (\log n)^{\frac{5}{2}}}{D^5 n^5 (5 \log n)^2} = C' \sqrt{\log n}.$$

## Extending the argument

- Let us consider

$$\eta = \lambda|\varphi_n\rangle\langle\varphi_n| + (1 - \lambda)\frac{1}{n^2}.$$

Note that

$$\eta^{\otimes k} = \lambda^k|\varphi_{n^k}\rangle\langle\varphi_{n^k}| + \dots\dots\dots$$

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- Since  $G_{n^k}$  has positive coefficients we know that

$$LV(\eta^{\otimes k}) \geq \frac{1}{\omega(G_{n^k})} \langle G_{n^k}, P \rangle \geq C \lambda^k \frac{n^k}{(\log n^k)^2} = C \frac{(\lambda n)^k}{k^2 (\log n)^2}.$$

# Consequences

- Superactivation of quantum nonlocality:

J. Barrett showed that there exist **local** isotropic states

$$\eta = \lambda |\varphi_n\rangle\langle\varphi_n| + (1 - \lambda) \frac{1}{n^2} \text{ for } \lambda > \frac{1}{n}. \text{ Therefore,}$$

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- Unbounded almost-superactivation of quantum nonlocality:

For every  $\epsilon > 0$  and  $\delta > 0$  there exists a high enough  $n$  such that  $\eta = \lambda_n|\varphi_n\rangle\langle\varphi_n| + (1 - \lambda_n)\frac{1}{n^2}$  verifies that

$$LV(\eta) \leq 1 + \epsilon \quad \text{and} \quad LV(\eta^{\otimes 5}) \geq \delta.$$



*Thank you very much!*