

# Quantum de Finetti theorems under local measurements

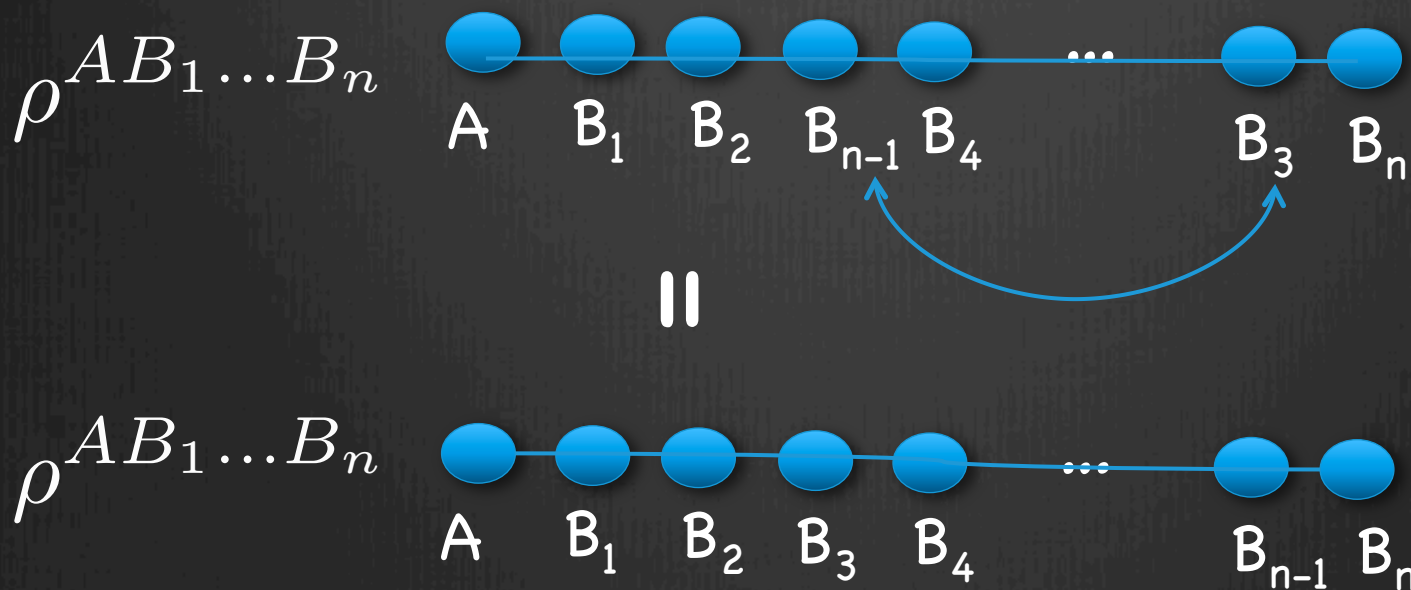
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based on [arXiv:1210.6367](https://arxiv.org/abs/1210.6367)  
joint work with Fernando Brandão (ETH)

# Symmetric States

$\rho^{AB_1 \dots B_n}$  is permutation symmetric in the B subsystems if for every permutation  $\pi$ ,

$$\rho^{AB_1 \dots B_n} = \rho^{AB_{\pi(1)} \dots B_{\pi(n)}}$$



# Quantum de Finetti Theorem

**Theorem** [Christandl, Koenig, Mitchison, Renner '06]

Given a state  $\rho^{AB_1 \dots B_n}$  symmetric under exchange of  $B_1 \dots B_n$ , there exists  $\mu$  such that

$$\left\| \rho^{AB_1 \dots B_k} - \int \mu(d\sigma) \rho_\sigma \otimes \sigma^{\otimes k} \right\|_1 \leq \frac{d^2 k}{n}$$

builds on work by [Størmer '69], [Hudson, Moody '76], [Raggio, Werner '89]  
[Caves, Fuchs, Sachs '01], [Koenig, Renner '05]

Proof idea:

Perform an informationally complete measurement of  $n-k$   $B$  systems.

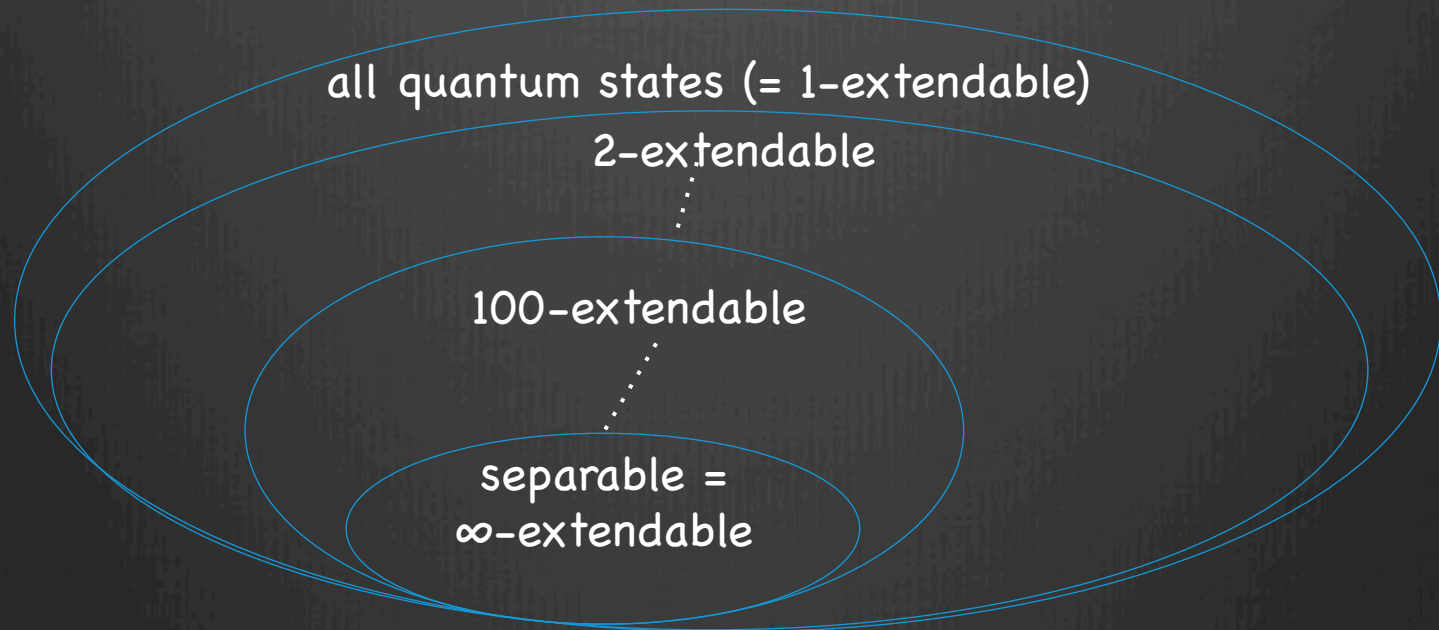
**Applications:**

information theory: tomography, QKD, hypothesis testing

algorithms: approximating separable states, mean-field theory

# Quantum de Finetti Theorem as Monogamy of Entanglement

Definition:  $\rho^{AB}$  is **n-extendable** if there exists an extension  $\rho^{AB_1 \dots B_n}$  with  $\rho^{AB} = \rho^{AB_i}$  for each  $i$ .



Algorithms: Can search/optimize over n-extendable states in time  $d^{O(n)}$ .

Question: How close are n-extendable states to separable states?

# Quantum de Finetti theorem

Theorem [Christandl, Koenig, Mitchison, Renner '06]

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Difficulty:

1. Parameters are, in many cases, **too weak**.
2. They are also essentially **tight**.

Way forward:

1. Change definitions (of error or i.i.d.)
2. Obtain better scaling

# relaxed/improved versions

Two examples known:

1. Exponential de Finetti Theorem: [Renner '07]

error term  $\exp(-\Omega(n-k))$ .

Target state convex combination of “almost i.i.d.” states.

2. measure error in 1-LOCC norm [Brandão, Christandl, Yard '10]

For error  $\varepsilon$  and  $k=1$ , requires  $n \sim \varepsilon^{-2} \log|A|$ .

This talk

improved de Finetti theorems for local  
measurements

# main idea

## use information theory

$$\log |A| \geq I(A:B_1 \dots B_n) = I(A:B_1) + I(A:B_2|B_1) + \dots + I(A:B_n|B_1 \dots B_{n-1})$$

repeatedly uses chain rule:  $I(A:BC) = I(A:B) + I(A:C|B)$

→  $I(A:B_t|B_1 \dots B_{t-1}) \leq \log(|A|)/n$  for some  $t \leq n$ .

If  $B_1 \dots B_n$  were classical, then we would have

$$\rho^{AB} = \rho^{AB_t} = \sum_i \pi_i \rho_i^{AB} \approx \text{separable}$$

Question:

How to make  $B_{1 \dots n}$  classical?

distribution  
on  $B_1 \dots B_{t-1}$

≈ product state  
(cf. Pinsker ineq.)



# Answer: measure!

Fix a measurement  $M:B \rightarrow Y$ .

$I(A:B_t|B_1 \dots B_{t-1}) \leq \varepsilon$  for the measured state  $(\text{id} \otimes M^{\otimes n})(\rho)$ .

Then

- $\rho^{AB}$  is hard to distinguish from  $\sigma \in \text{Sep}$  if we first apply  $(\text{id} \otimes M)$
- $\|(\text{id} \otimes M)(\rho - \sigma)\| \leq \text{small}$  for some  $\sigma \in \text{Sep}$ .

## Theorem

Given a state  $\rho^{AB_1 \dots B_n}$  symmetric under exchange of  $B_1 \dots B_n$ ,  
and  $\{\Lambda_i\}$  a collection of operations from  $A \rightarrow X$ ,

$$\min_{\sigma \in \text{Sep}} \max_M \mathbb{E}_i \left\| (\Lambda_i^A \otimes M^B)(\rho^{AB} - \sigma^{AB}) \right\|_1 \leq \sqrt{\frac{2 \ln |X|}{n}}$$

Cor: setting  $\Lambda = \text{id}$  recovers [Brandão, Christandl, Yard '10] 1-LOCC result.



# advantages/extensions

## Theorem

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1. Simpler proof and better constants
2. Bound depends on  $|X|$  instead of  $|A|$  (can be  $\infty$  dim)
3. Applies to general non-signalling distributions
4. There is a multipartite version (multiply error by  $k$ )
5. Efficient “rounding” (i.e.  $\sigma$  is explicit)
6. Symmetry isn't required (see Fernando's talk on Thursday)

# applications

- **nonlocal games**

Adding symmetric provers “immunizes” against entanglement / non-signalling boxes. (Caveat: needs uncorrelated questions.)  
Conjectured improvement would yield NP-hardness for 4 players.

- **BellQMA(poly) = QMA**

Proves Chen-Drucker  $\text{SAT} \in \text{BellQMA}_{\log(n)}(\sqrt{n})$  protocol is optimal.

- **pretty good tomography** [Aaronson '06]

on permutation-symmetric states (instead of product states)

- **convergence of Lasserre hierarchy** for polynomial optimization  
see also 1205.4484 for connections to small-set expansion

# open questions

- Is  $\text{QMA}(2) = \text{QMA}$ ? Is  $\text{SAT} \in \text{QMA}_{\sqrt{n}}(2)_{1,1/2}$  optimal? (Would follow from replacing 1-LOCC with SEP-YES.)

- Can we reorder our quantifiers to obtain

$$\min_{\sigma \in \text{Sep}} \mathbb{E}_i \max_M \left\| (\Lambda_i^A \otimes M^B)(\rho^{AB} - \sigma^{AB}) \right\|_1 \leq \sqrt{\frac{2 \ln |X|}{n}}?$$

(no-signalling analogue is FALSE assuming  $P \neq NP$ )

- The usual de Finetti questions:
  - better counter-examples
  - how much does it help to add PPT constraints?