

A fault-tolerant one-way quantum computer

Robert Raussendorf¹, Jim Harrington² and
Kovid Goyal¹

1: California Institute of Technology,

2: Los Alamos National Laboratory

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Main idea

Universal quantum computation by local measurements:

- A three-dimensional cluster state is a *fault-tolerant fabric*.

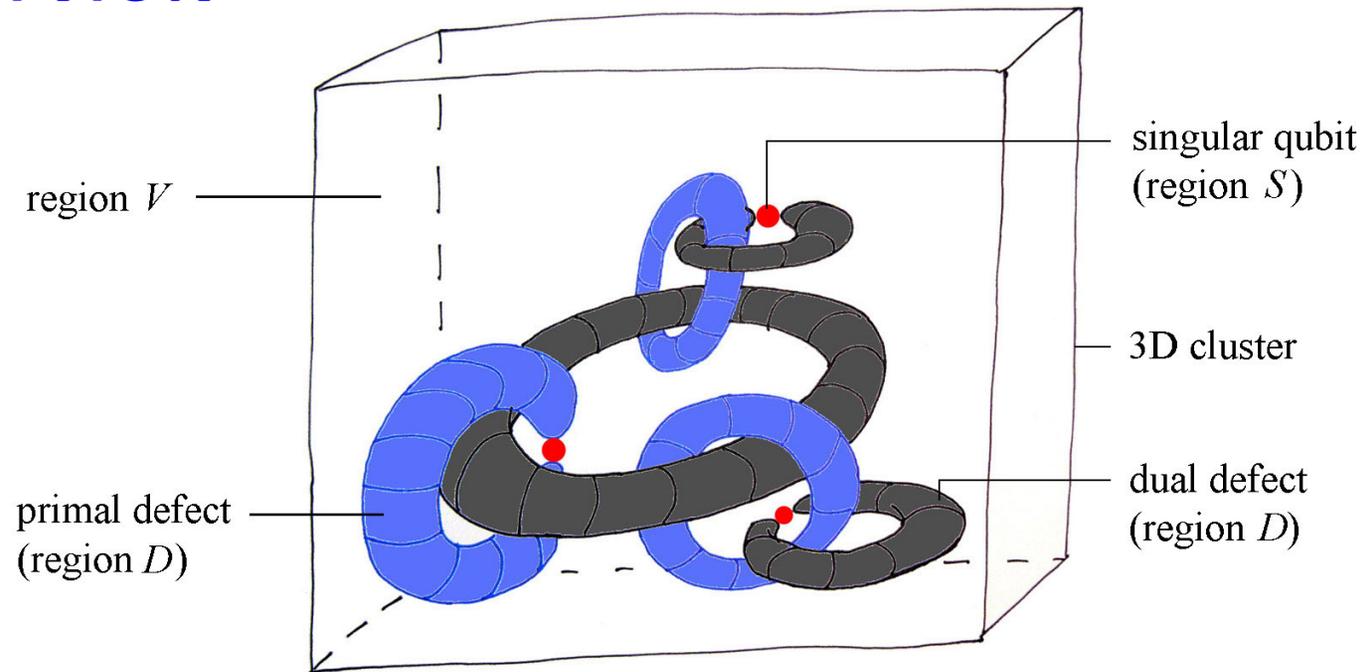


Make use of geometry!

Result

- Geometric constraint: only *local* and *next-neighbor* interaction of qubits on a three-dimensional lattice required.
- Fault-tolerance threshold of 1.1×10^{-3} (each source).
Error sources: preparation, gates, storage and measurement.
- Polynomial overhead.

Overview



Three cluster regions:

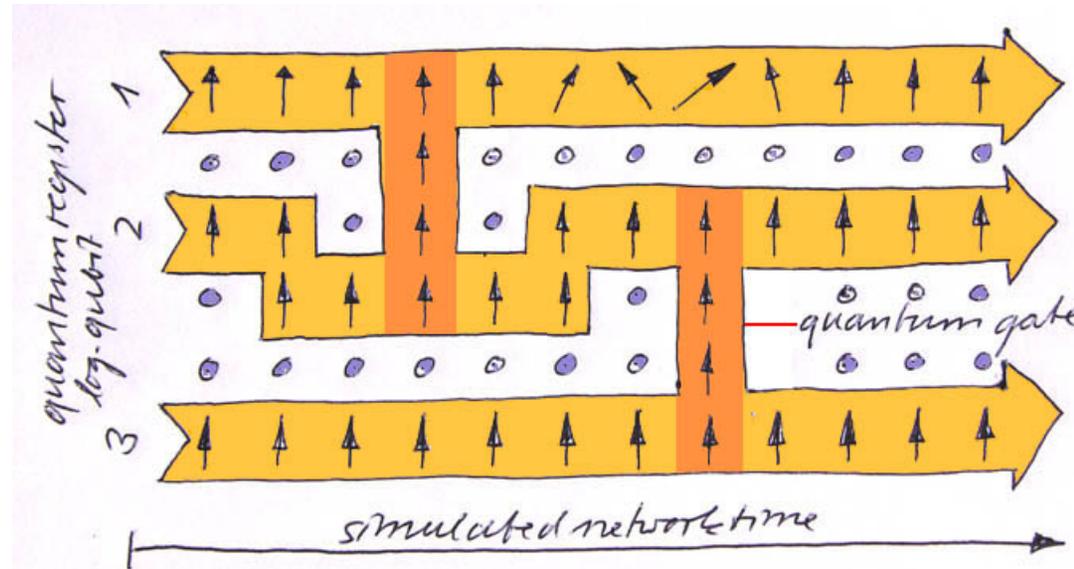
V (Vacuum), D (Defect) and S (Singular qubits).

Qubits $q \in V$: local X -measurements,
Qubits $q \in D$: local Z -measurements,
Qubits $q \in S$: local measurements of $\frac{X \pm Y}{\sqrt{2}}$.

Preliminaries:

The one-way quantum computer, cluster states and topological codes

The one-way quantum computer



measurement of Z (\odot), X (\uparrow), $\cos \alpha X + \sin \alpha Y$ (\nearrow)

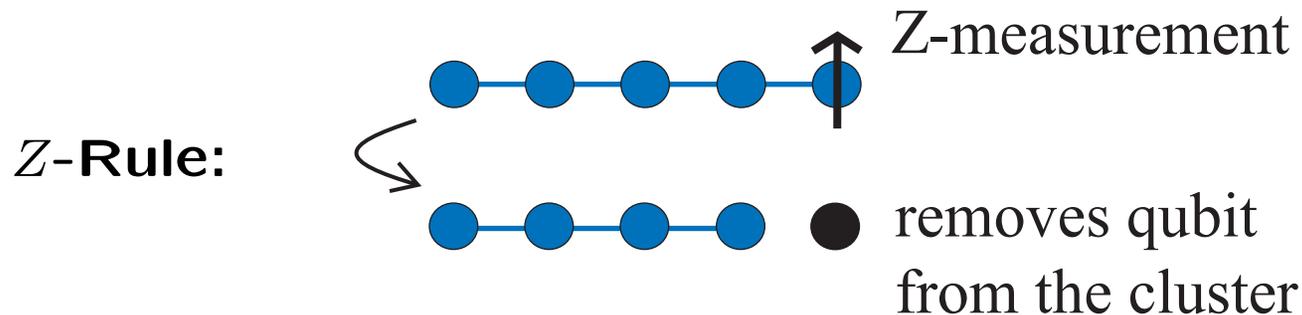
- Universal computational resource: cluster state.
- Information written onto the cluster, processed and read out by one-qubit measurements only.

Cluster states

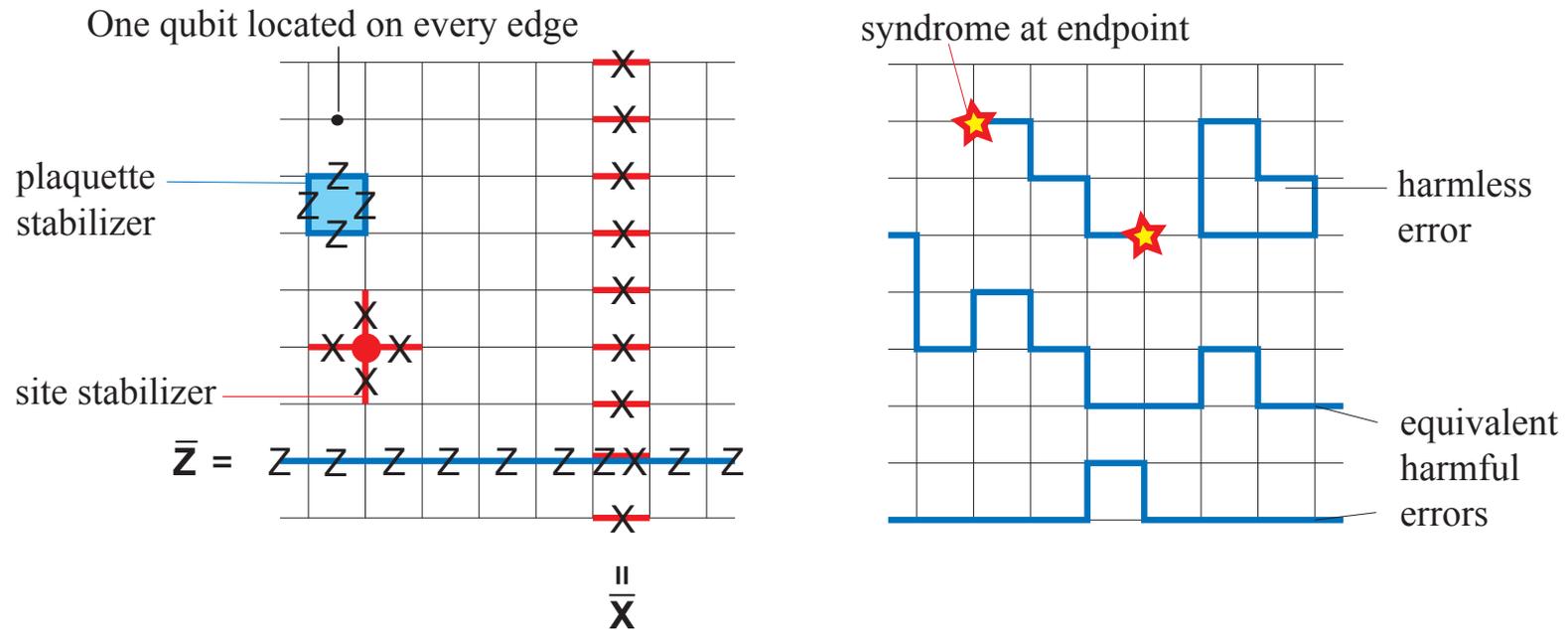
A cluster state $|\phi\rangle_{\mathcal{C}}$ on a cluster \mathcal{C} is the single common eigenstate of the stabilizer operators

$$K_a = X_a \bigotimes_{b \in N(a)} Z_b, \quad \forall a \in \mathcal{C}, \quad (1)$$

where $b \in N(a)$ if a, b are spatial next neighbors in \mathcal{C} .



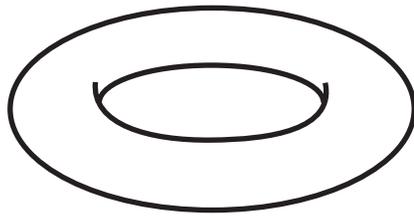
Topological quantum codes



- Errors are represented by chains.
- Homologically equivalent chains correspond to physically equivalent errors.
- Harmful errors stretch across the entire lattice (rare events).

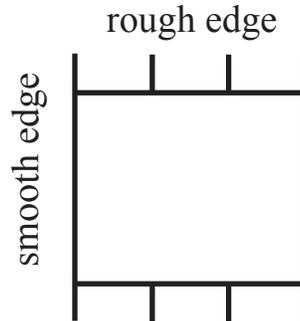
A. Kitaev, quant-ph/9707021 (1997).

Topological quantum codes



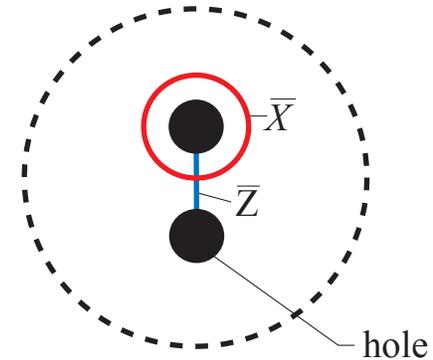
Torus

2 Qubits



Plane segment

1 Qubit

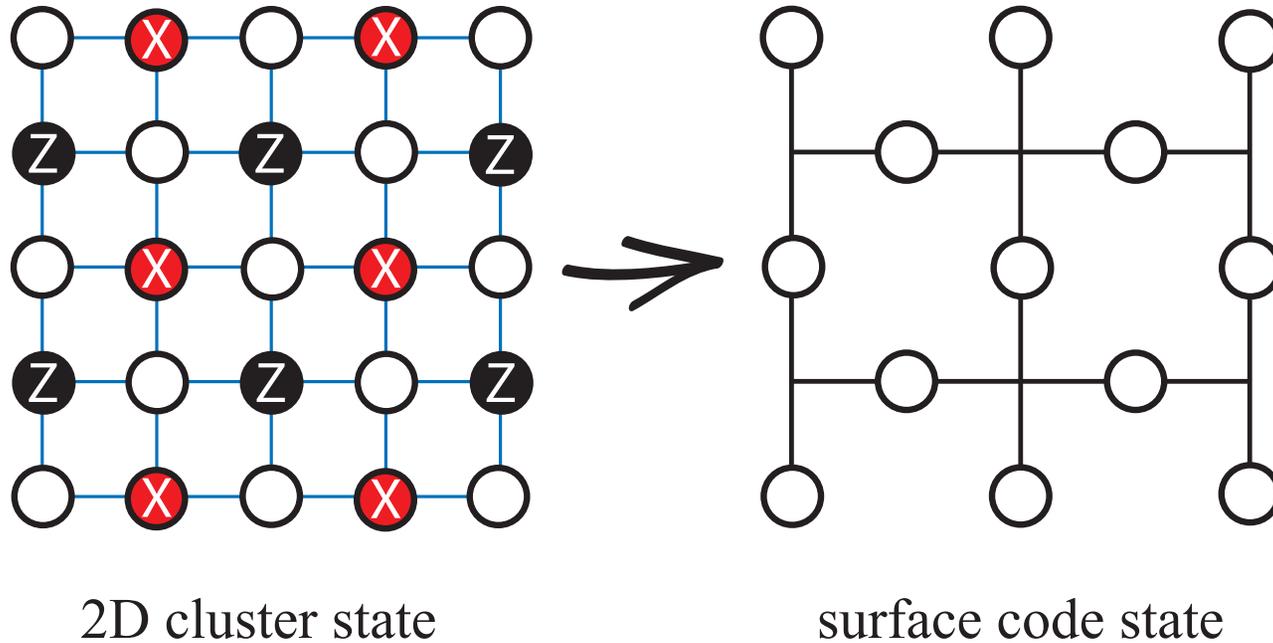


Plane with 2 holes

1 Qubit

- Storage capacity of the code depends upon the topology of the code surface.

Link



- Obtain surface code state from 2D cluster state via regular pattern of Z - and X -measurements.

Talk outline

Methods

Error-correction
in 3D cluster states

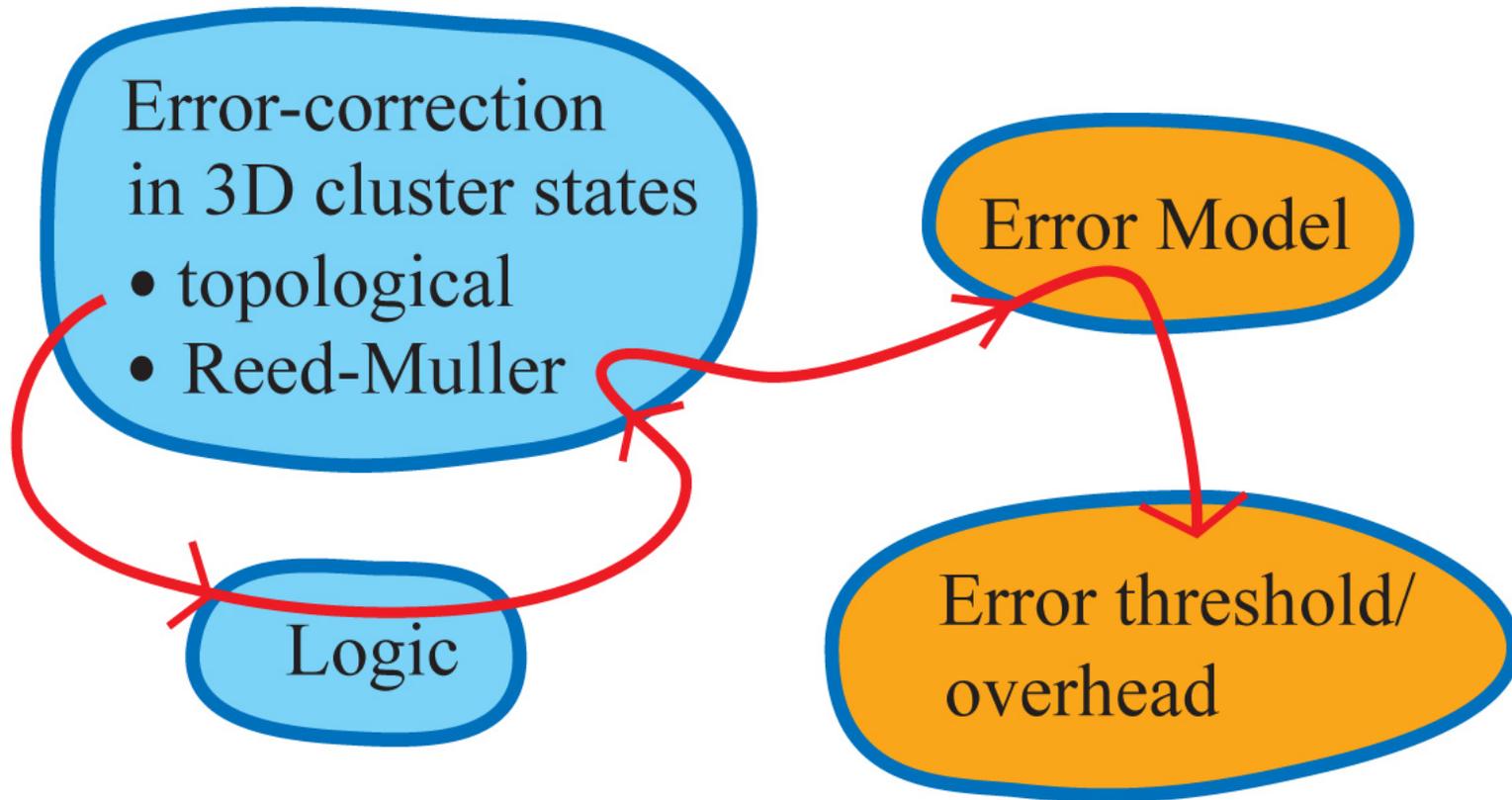
- topological
- Reed-Muller

Logic

Numbers

Error Model

Error threshold/
overhead

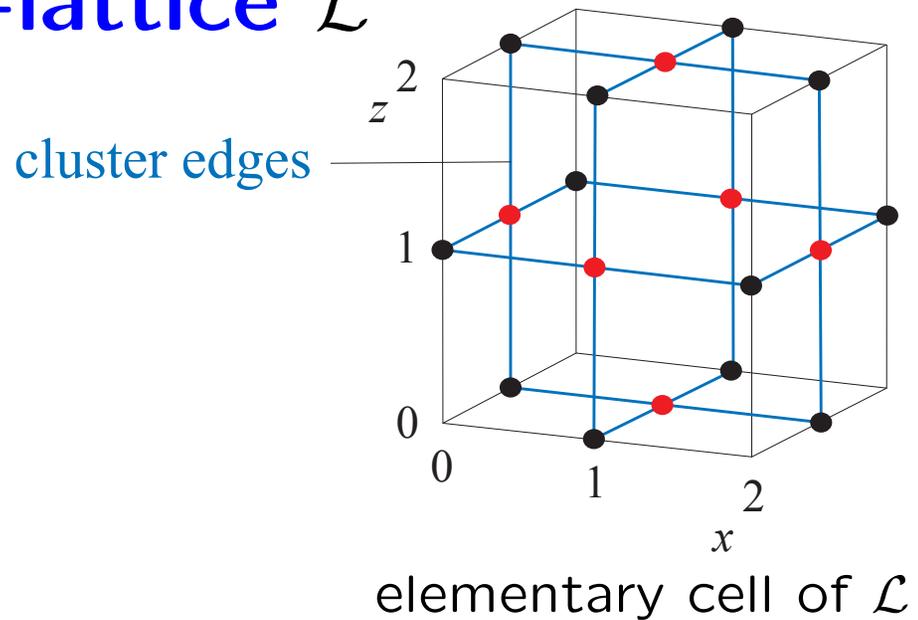


Part I:

Error correction in 3D cluster states

Cluster states in three spatial dimensions provide intrinsic topological error correction.

Cluster \mathcal{C} and bcc-lattice \mathcal{L}



- qubit location: (even, odd, odd) - face of \mathcal{L} ,
- qubit location: (odd, odd, even) - edge of \mathcal{L} ,
- syndrome location: (odd, odd, odd) - cube of \mathcal{L} ,
- syndrome location: (even, even, even) - site of \mathcal{L} .

Topological error correction in V

Measurement pattern:

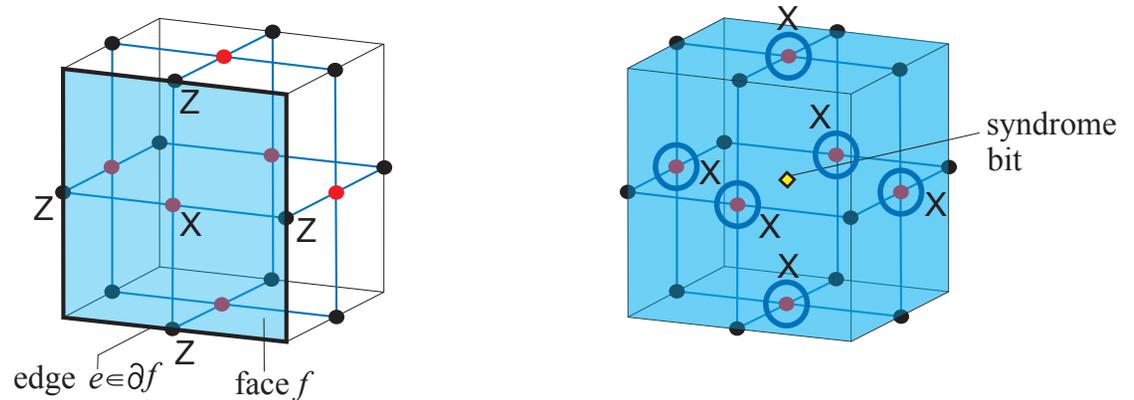
- The qubits $q \in V$ are individually measured in the X -basis.

Errors:

- Consider probabilistic Pauli errors.
- Sufficient to consider Z -errors.

(X -errors are absorbed into the X -measurement, $\frac{I \pm X}{2} X = \pm \frac{I \pm X}{2}$.)

Homology



- Stabilizer elements associated with faces f of \mathcal{L} :

$$K(f) = \bigotimes_{a \in f} X_a \bigotimes_{b \in \partial f} Z_b. \quad (2)$$

- Stabilizer for syndrome ($[K(f), X_q] = 0 \forall q \in V$):

$$\partial f = 0. \quad (3)$$

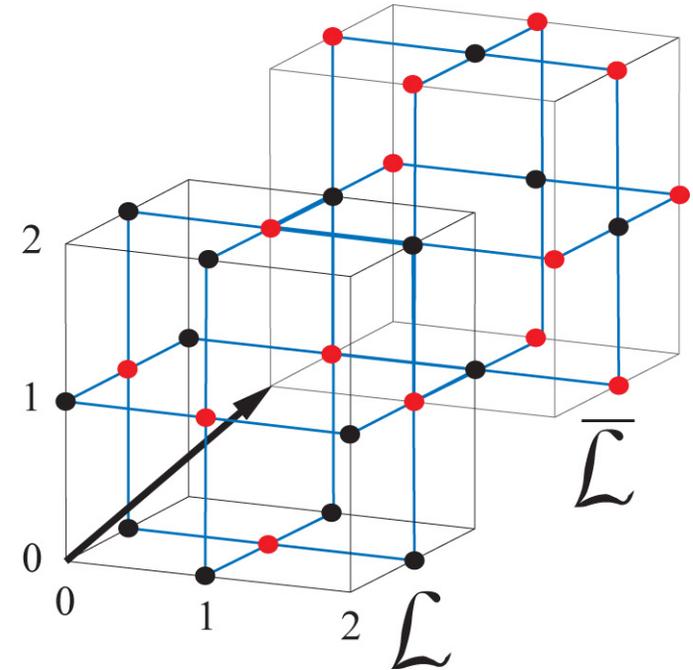
- One syndrome bit per cell of \mathcal{L} . Protects the face qubits.

What about the edge qubits?

Lattice duality $\mathcal{L} \longleftrightarrow \bar{\mathcal{L}}$

Translation by vector $(1, 1, 1)^T$:

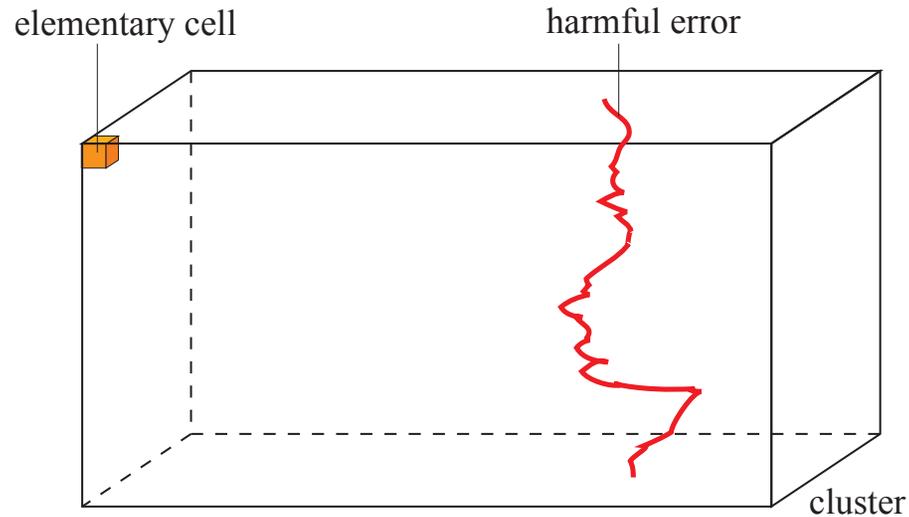
- Cluster \mathcal{C} invariant,
- \mathcal{L} (primal) $\longrightarrow \bar{\mathcal{L}}$ (dual).



$$\begin{aligned}
 \text{face of } \mathcal{L} & - \text{edge of } \bar{\mathcal{L}}, \\
 \text{edge of } \mathcal{L} & - \text{face of } \bar{\mathcal{L}}, \\
 \text{site of } \mathcal{L} & - \text{cube of } \bar{\mathcal{L}}, \\
 \text{cube of } \mathcal{L} & - \text{site of } \bar{\mathcal{L}},
 \end{aligned}
 \tag{4}$$

- Edge qubits protected by stabilizer on dual lattice $\bar{\mathcal{L}}$.
- Many objects in this scheme exist as 'primal' and 'dual'.

Topological error correction in V

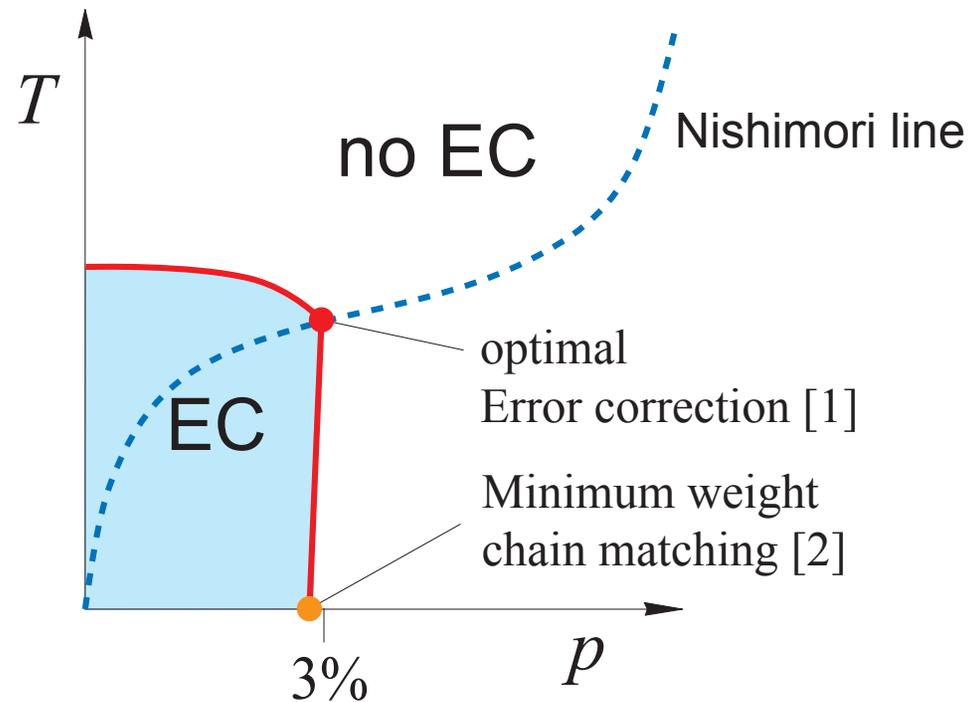


- One syndrome bit for each elementary cell of \mathcal{L} .
 - Harmful errors stretch across entire lattice \mathcal{L} .
- > Leads to *Random plaquette Z_2 -gauge model* (RPGM) [1].

[1] Dennis et al., quant-ph/0110143 (2001).

RPGM: schematic phase diagram

Map error correction to statistical mechanics:



[1] T. Ohno et al., quant-ph/0401101 (2004).

[2] E. Dennis et al., quant-ph/0110143 (2001); J. Edmonds, Canadian J. Math. **17**, 449 (1965).

Cluster region V



Defects $d \in D$

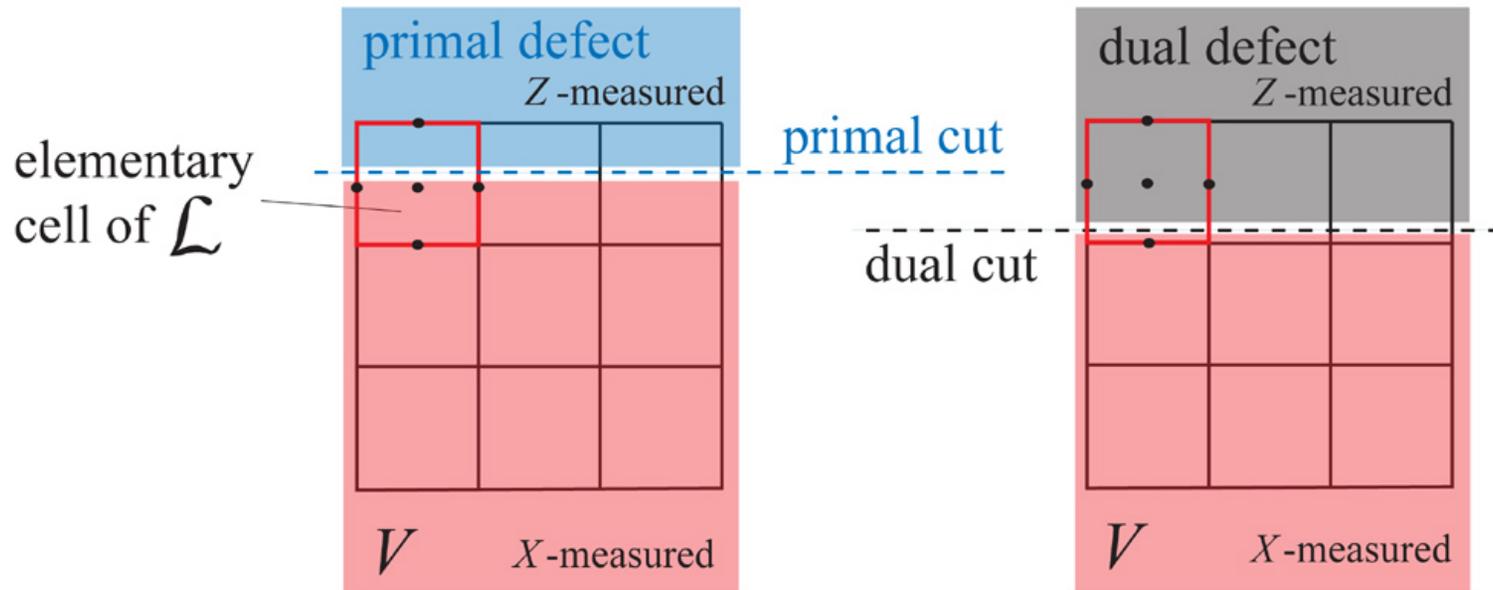


Singular qubits

Part II: Quantum Logic

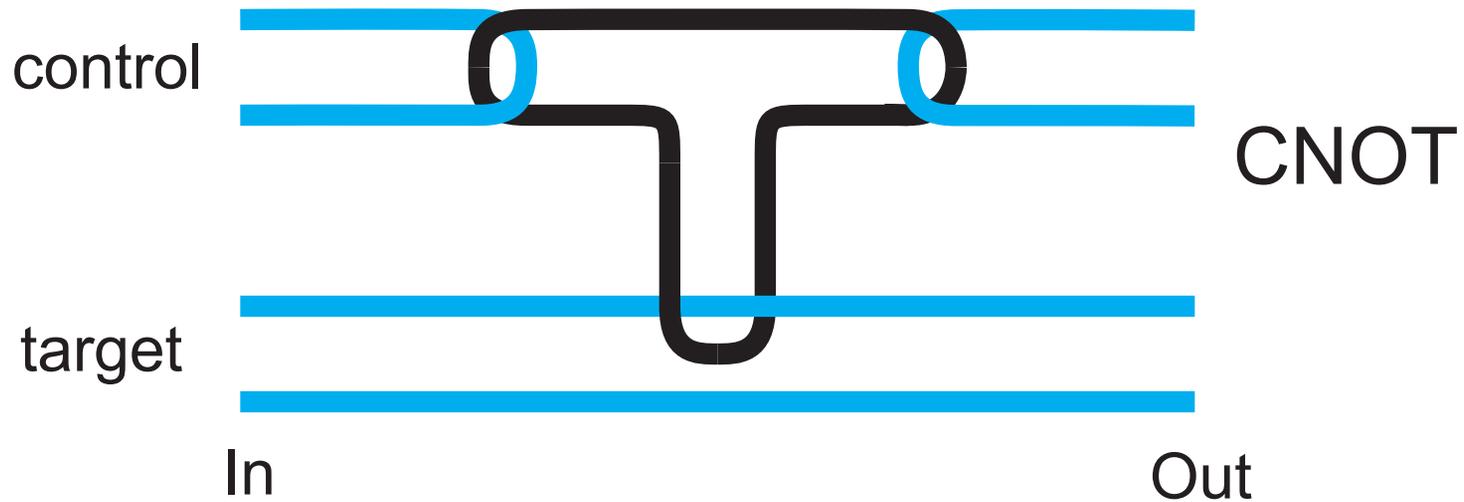
Fault-tolerant quantum logic is realized via topologically entangled engineered lattice defects.

Defects



- Defects are regions of the cluster where qubits are measured in the Z -basis.
- Defects create cluster boundaries (cuts).
- There are *primal* and *dual* defects.

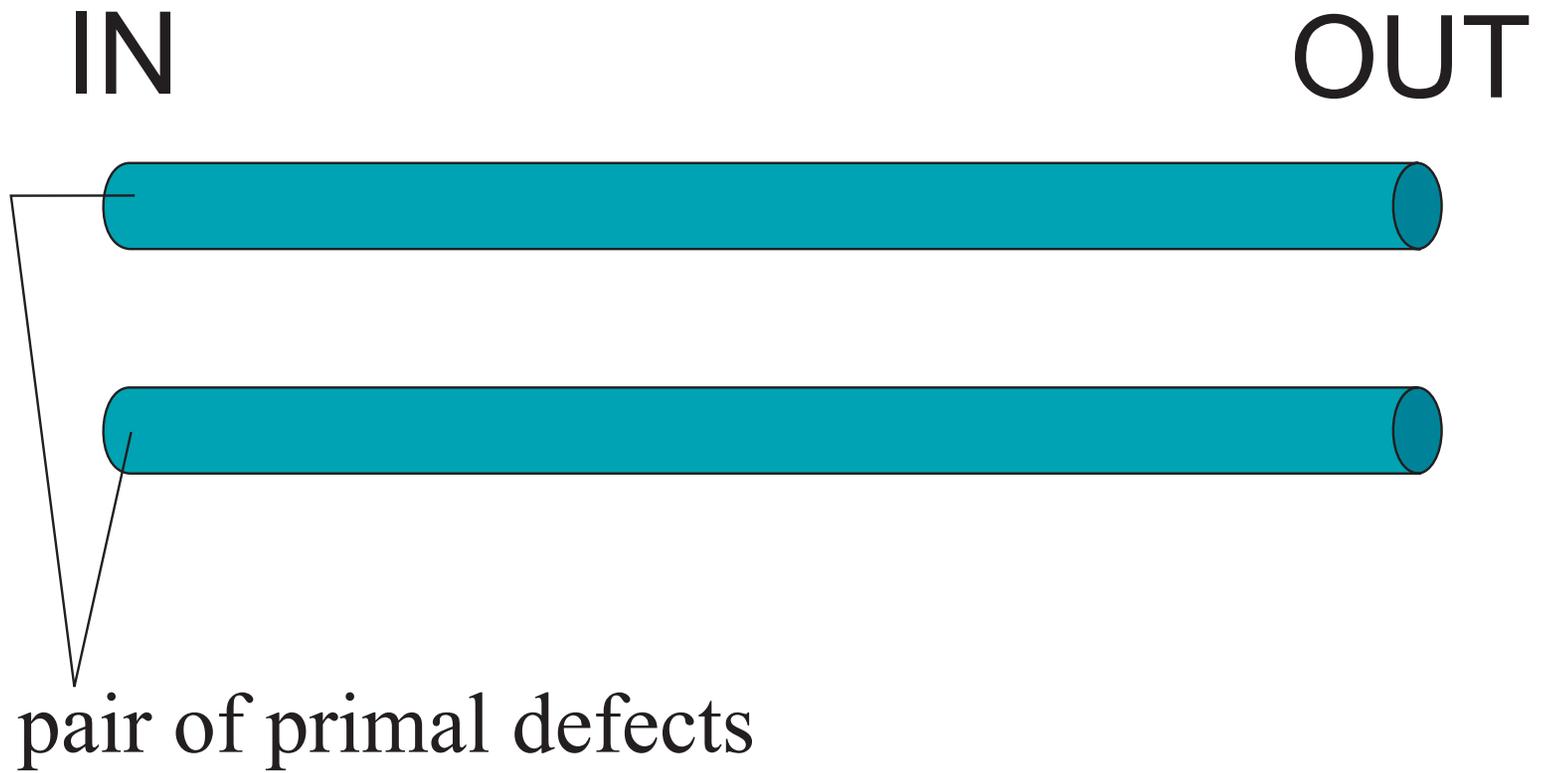
Defects for quantum logic



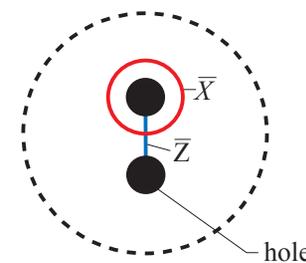
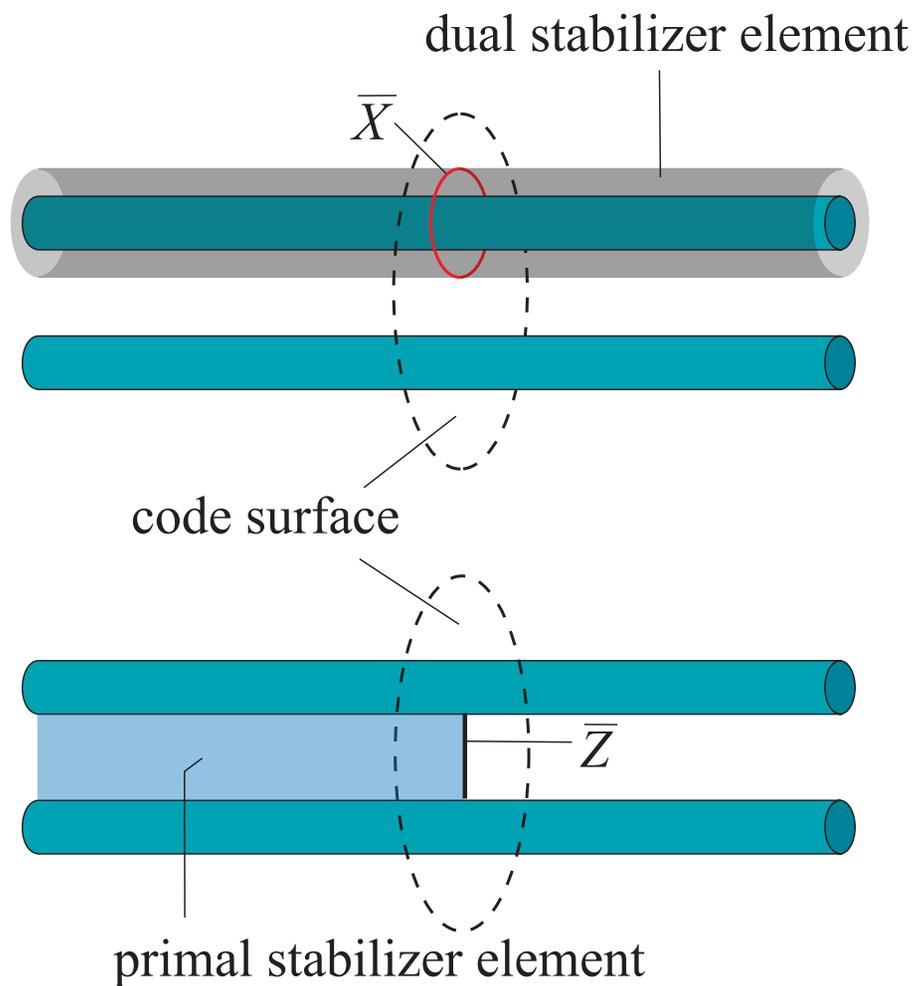
*A quantum circuit is encoded in the way
primal and dual defects are wound around another.*

Quantum gates, Part I

Piece of wire



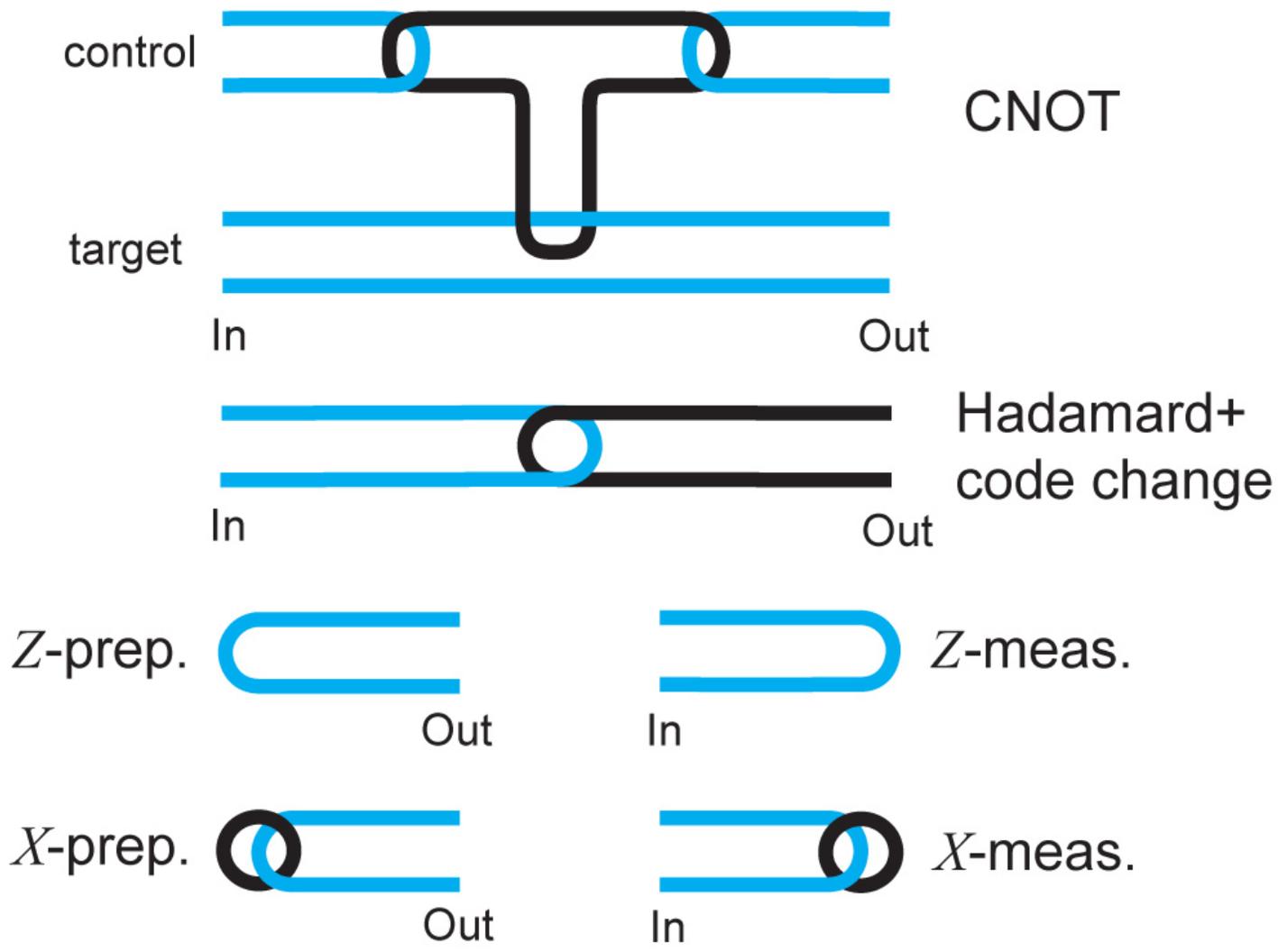
Quantum gates, Part I



Plane with 2 holes

1 Qubit

Quantum gates, Part I



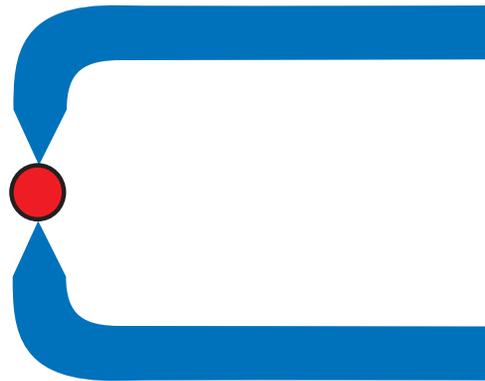
- Displayed fault-tolerant gates are not universal.
- Need one non-Clifford element:
fault-tolerant measurement of $\frac{X \pm Y}{\sqrt{2}}$.

Cluster region V	✓
Defects $d \in D$	✓
Singular qubits	←

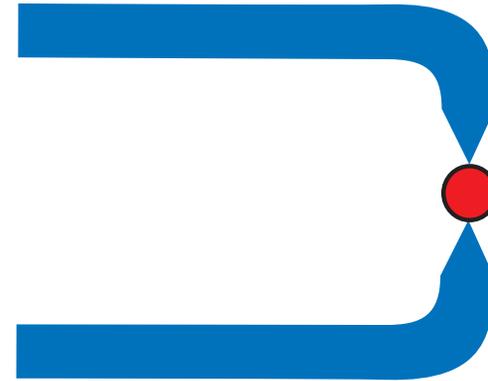
Quantum gates, Part II

Encoder and decoder for surface code:

singular
qubit



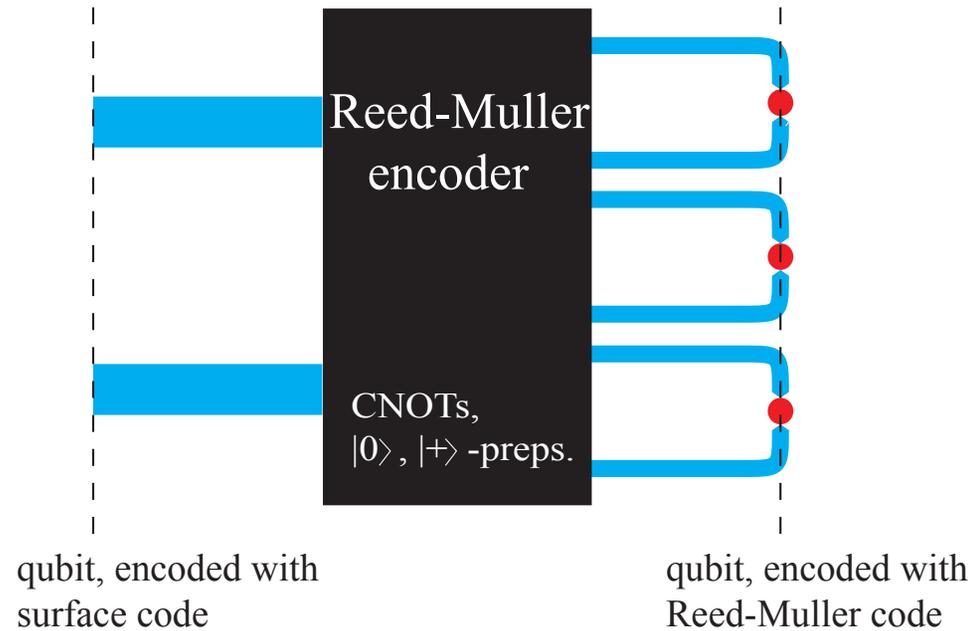
Encoder



Decoder

Quantum gates, Part II

A circuit for code-conversion:



- Reed-Muller code: Fault-tolerant measurement of $\frac{\bar{X} \pm \bar{Y}}{\sqrt{2}}$ via *local* measurements of $\frac{X_a \pm Y_a}{\sqrt{2}}$ and classical post-processing.

-> *Fault-tolerant universal gate set complete.*

Part III:
The Error Model

Error model:

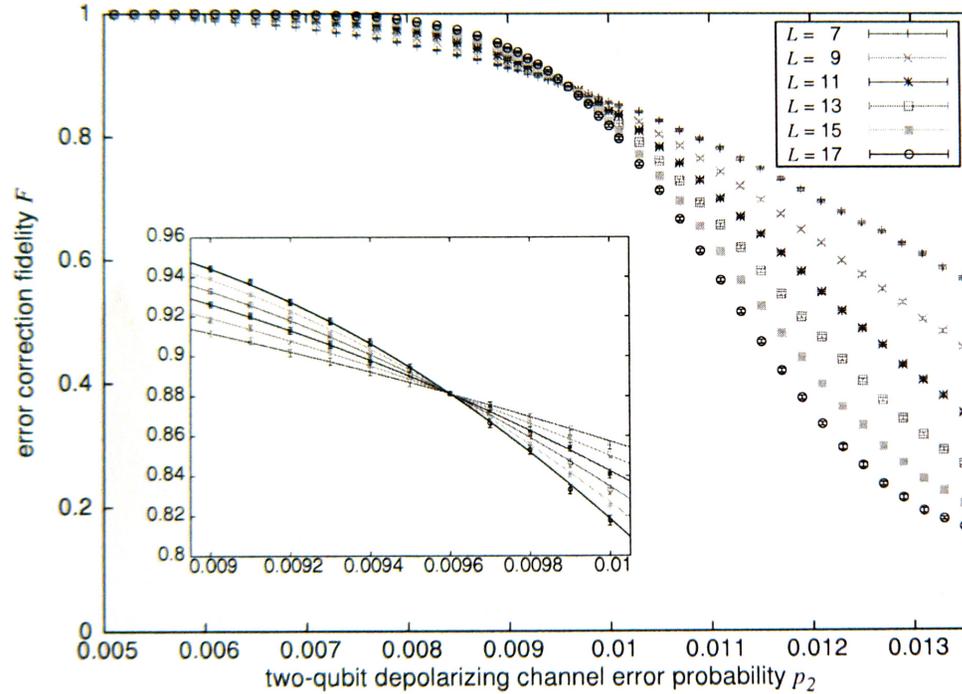
- Cluster state created in a 4-step sequence of $\Lambda(Z)$ -gates from product state $\bigotimes_{a \in \mathcal{C}} |+\rangle_a$.
- Error sources:
 - **$|+\rangle$ -preparation**: Perfect preparation followed by 1-qubit partially depolarizing noise with probability p_P .
 - **$\Lambda(Z)$ -gates**: Perfect gates followed by 2-qubit partially depolarizing noise with probability p_2 .
 - **Memory**: 1-qubit partially depolarizing noise with probability p_S per time step.
 - **Measurement**: Perfect measurement preceded by 1-qubit partially depolarizing noise with probability p_M .
- 3D cluster state created in slices of fixed thickness.
- Instant classical processing.

Part IV:

Threshold and overhead

The fault-tolerance threshold is 1.1×10^{-3} for each source. The overhead is polynomial.

Topological error-correction in V



$$\begin{aligned}
 p_{2,c} &= 9.6 \times 10^{-3}, & \text{for } p_P = p_S = p_M = 0, \\
 p_c &= 5.8 \times 10^{-3}, & \text{for } p_P = p_S = p_M = p_2 =: p.
 \end{aligned} \tag{5}$$

Reed-Muller error-correction in S

Error budget from Reed-Muller concatenation threshold:

$$\frac{76}{15}p_2 + \frac{2}{3}p_P + \frac{4}{3}p_M + \frac{4}{3}p_S < \frac{1}{105}. \quad (6)$$

Specific parameter choices:

$$\begin{aligned} p_{2,c} &= 2.9 \times 10^{-3}, & \text{for } p_P = p_S = p_M = 0, \\ p_c &= 1.1 \times 10^{-3}, & \text{for } p_P = p_S = p_M = p_2 =: p. \end{aligned} \quad (7)$$

The Reed-Muller code sets the overall threshold.

Overhead

N : Number of non-Clifford operations in bare computation.

N_{ft} : Number of operations for fault-tolerant computation.

$$N_{\text{ft}} \leq N^2 (\log N)^{10.8}. \quad (8)$$

- Overhead is polynomial.
- Exponents may be reduced in more resourceful adaptations.

Summary

[quant-ph/0510135]

Scenario:

- Local and next-neighbor gates in 3D.

Numbers:

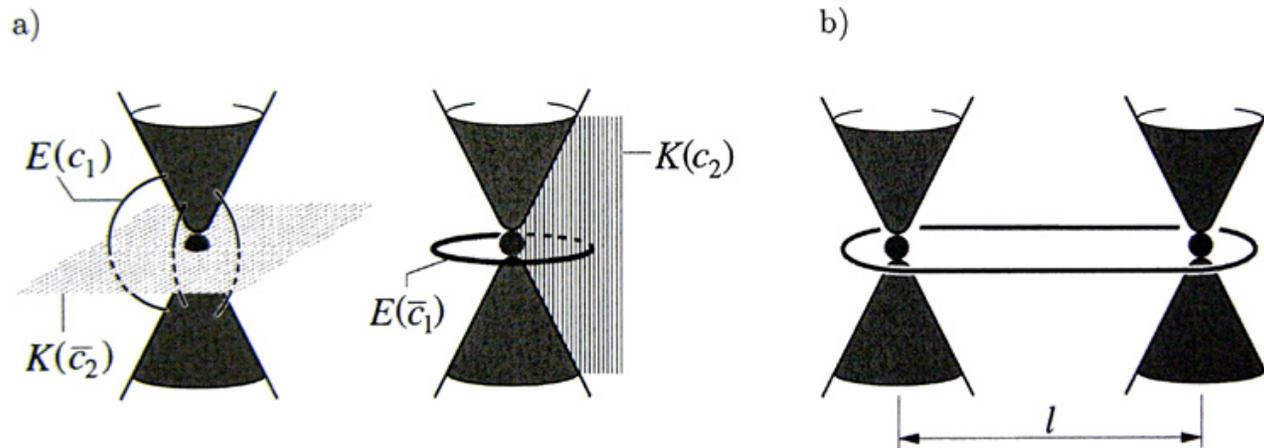
- Fault-tolerance threshold of 1.1×10^{-3} for preparation-, gate-, storage- and measurement error (each source).

Methods:

- Cluster states in three spatial dimensions provide intrinsic topological error correction related to the *Random plaquette Z_2 -gauge model*.
- Quantum logic is realized by topologically entangled *engineered lattice defects*.

Supplementary material

Local residual error on S -qubits



- Topological error correction breaks down near the S -qubits.
- Leads to *local* effective error on S -qubits.

The CNOT-gate

